

# Self-Stabilizing Leader Election in Polynomial Steps<sup>1</sup>

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**Anaïs Durand**   Franck Petit

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<sup>1</sup>This work has been partially supported by the LabEx PERSYVAL-Lab (ANR-11-LABX-0025-01) and the AGIR project DIAMS.

# Problem

- **Silent Self-stabilizing Leader Election**
- Model:
  - ▶ **Locally shared memory** model
  - ▶ Read/write atomicity
  - ▶ Distributed **unfair** daemon
- Network:
  - ▶ Any connected topology
  - ▶ Bidirectional
  - ▶ **Identified**
- **No global knowledge** on the network

# State of the Art

Model	Paper	Knowledge			Daemon	Complexity			Silent
		$D$	$N$	$B$		Memory	Rounds	Steps	
Message Passing	Afek, Bremler, 1998			x		$\Theta(\log n)$	$O(n)$	?	✓
	Awerbuch <i>et al</i> , 1993	x				$\Theta(\log D \log n)$	$O(D)$	?	✓
	Burman, Kutten, 2007	x				$\Theta(\log D \log n)$	$O(D)$	?	✓
Locally Shared Memory	Dolev, Herman, 1997		x		Fair	$\Theta(N \log N)$	$O(D)$	?	
	Arora, Gouda, 1994	x			Weakly Fair	$\Theta(\log N)$	$O(N)$	?	✓
	Datta <i>et al</i> , 2010				Unfair	unbounded	$O(n)$	?	✓
	Kravchik, Kutten, 2013				Synchronous	$\Theta(\log n)$	$O(D)$	?	✓
	Datta <i>et al</i> , 2011				Unfair	$\Theta(\log n)$	$O(n)$	?	✓

$D$ : Diameter

$D \geq \mathcal{D}$ : Upper bound on the diameter

$n$ : Number of nodes

$N \geq n$ : Upper bound on the number of nodes

$B$ : Upper bound on the link-capacity

# Our Contribution

## Algorithm $\mathcal{LE}$

- Memory requirement asymptotically optimal:  $\Theta(\log n)$  bits/process
- Stabilization time (worst case):
  - ▶  $3n + \mathcal{D}$  rounds
  - ▶ Lower Bound:  $\frac{n^3}{6} + \frac{5}{2}n^2 - \frac{11}{3}n + 2$  steps,
  - Upper Bound:  $\frac{n^3}{2} + 2n^2 + \frac{n}{2} + 1$  steps

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## Analytical Study of Datta *et al*, 2011<sup>2</sup>

- Stabilization time **not polynomial** in steps:
  - ▶  $\forall \alpha \geq 3, \exists$  networks and executions in  $\Omega(n^{\alpha+1})$  steps.

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<sup>2</sup>Datta, Larmore, and Vemula. Self-stabilizing Leader Election in Optimal Space under an Arbitrary Scheduler. 2011

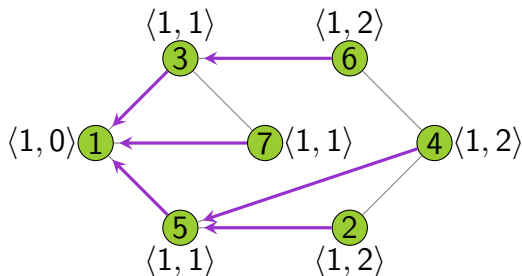
# Design of the Leader Election Algorithm

# Simplified Algorithm (Non Self-stabilizing)

Join a Tree

3 variables per process  $p$

- $p.idR \in \mathbb{N}$ : ID of the root
- $p.par \in \mathcal{N}_p \cup \{p\}$ : Parent pointer
- $p.level \in \mathbb{N}$ : Level



Key:  $\langle idR, level \rangle$

# Simplified Algorithm (Non Self-stabilizing)

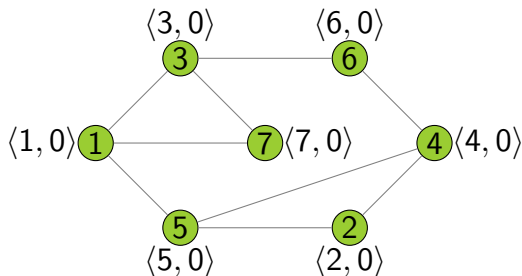
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### Initial Configuration

- $p.idR = p$
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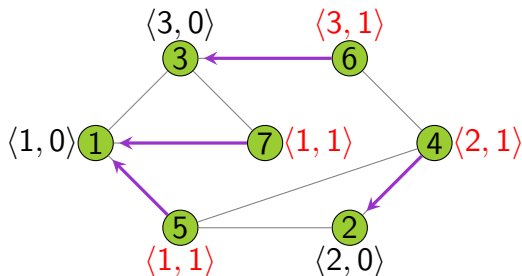
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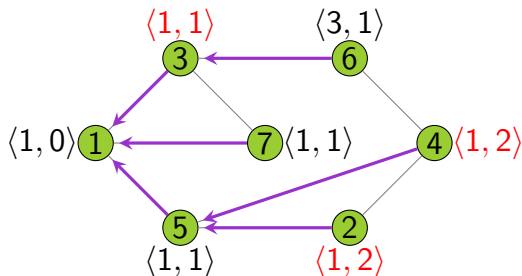
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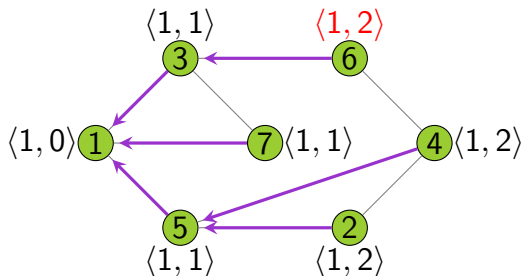
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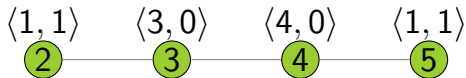
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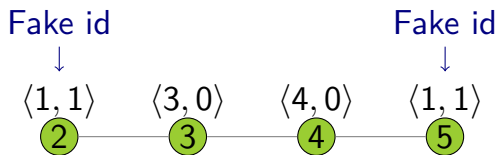
Self-stabilization  $\implies$  Arbitrary initialization



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# Simplified Algorithm (Non Self-Stabilizing)

Self-stabilization  $\implies$  Arbitrary initialization  $\implies$  Fake ids



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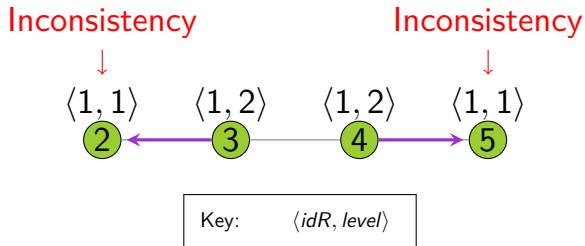
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# Simplified Algorithm: Removal of Fake Ids

Reset

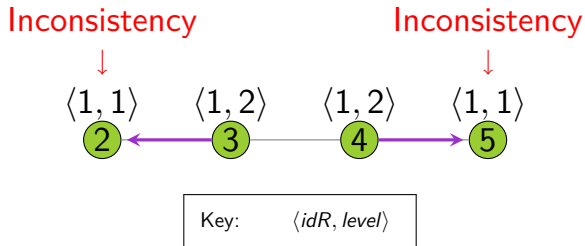


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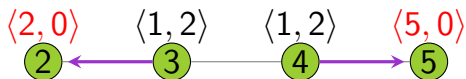


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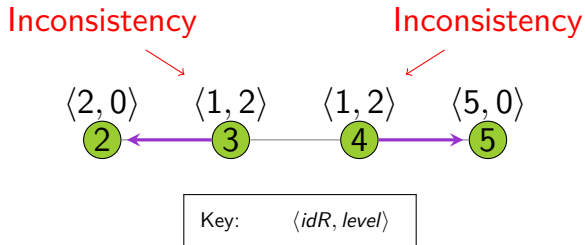
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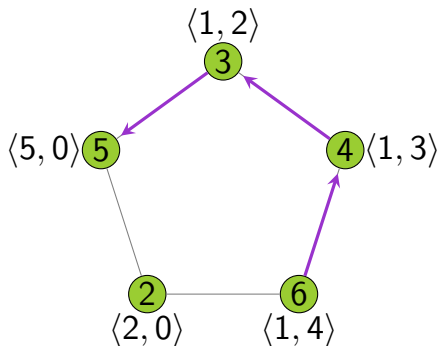
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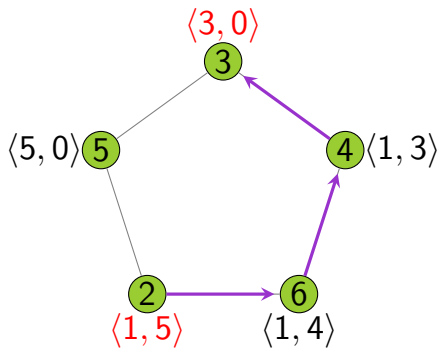
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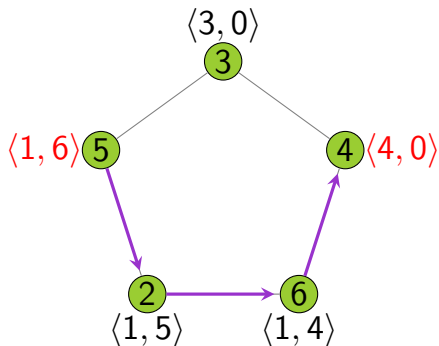
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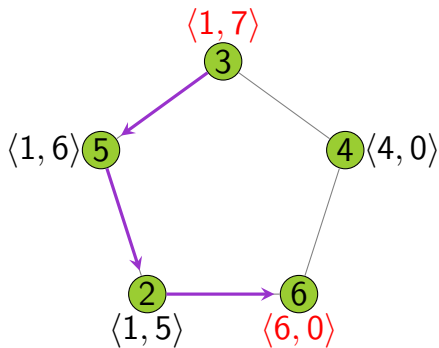
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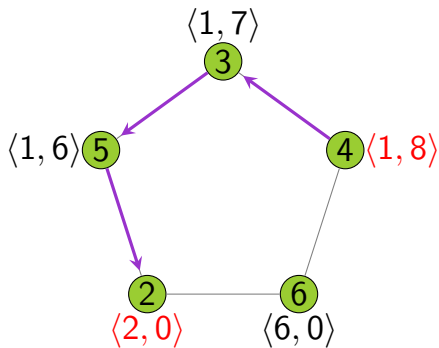
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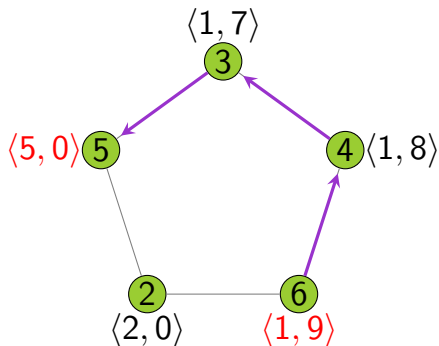
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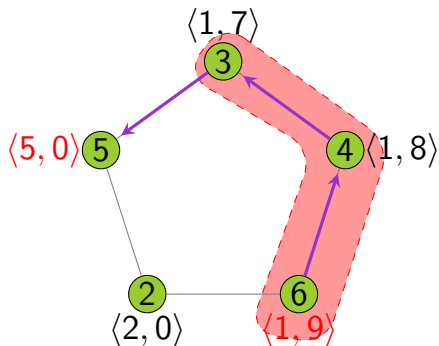
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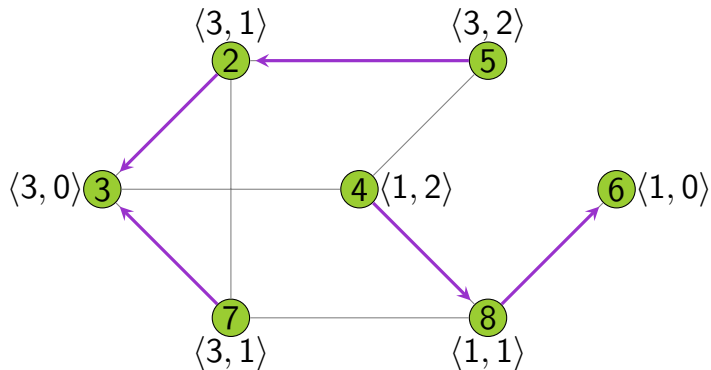
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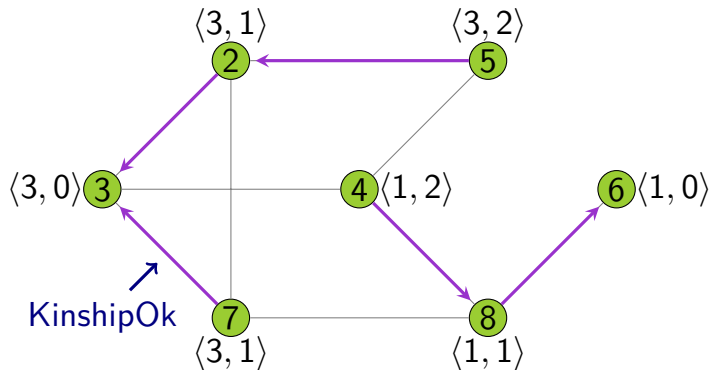
Key:  $\langle idR, level \rangle$

# Abnormal Trees



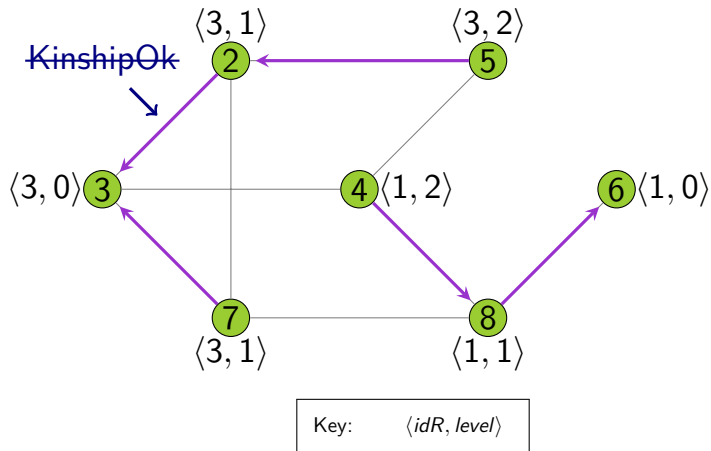
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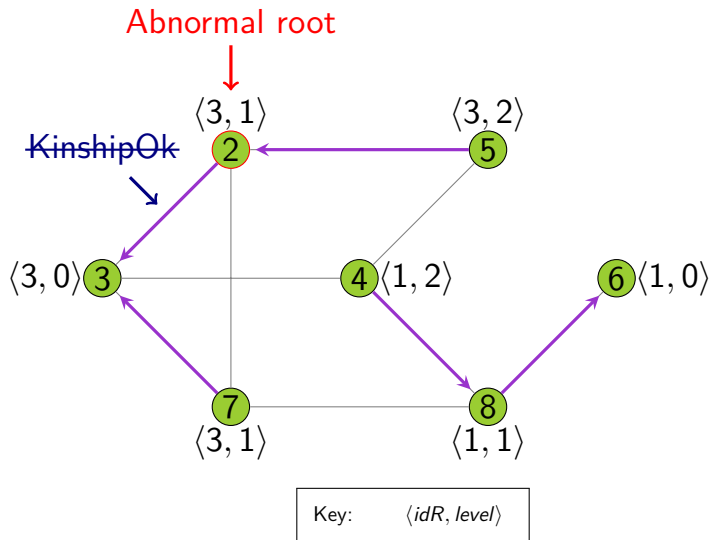


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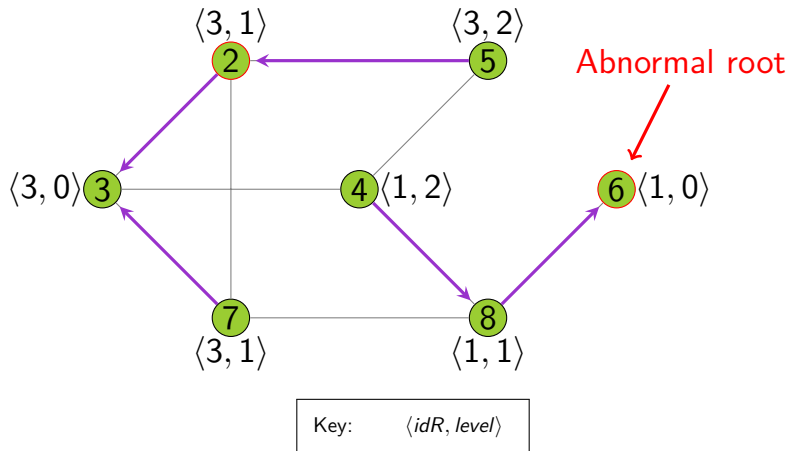
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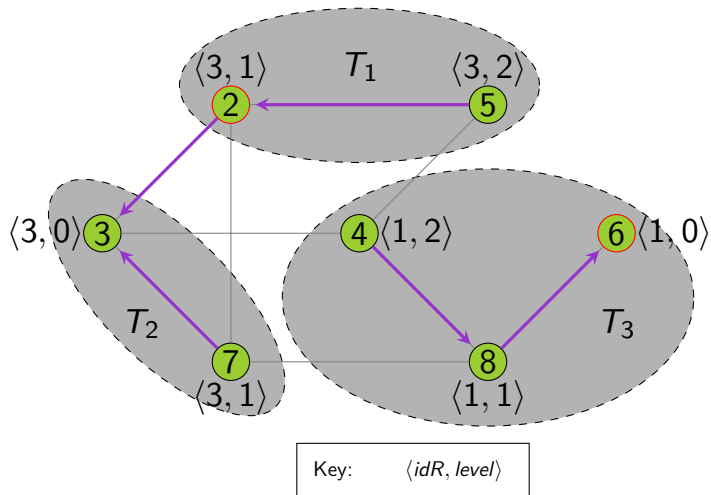
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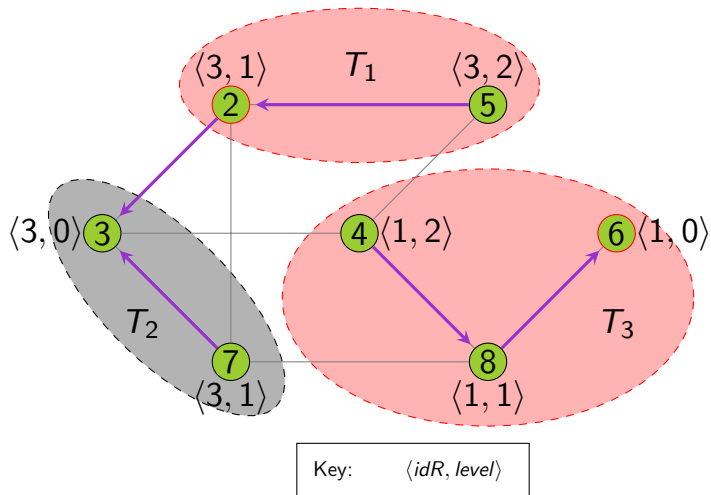


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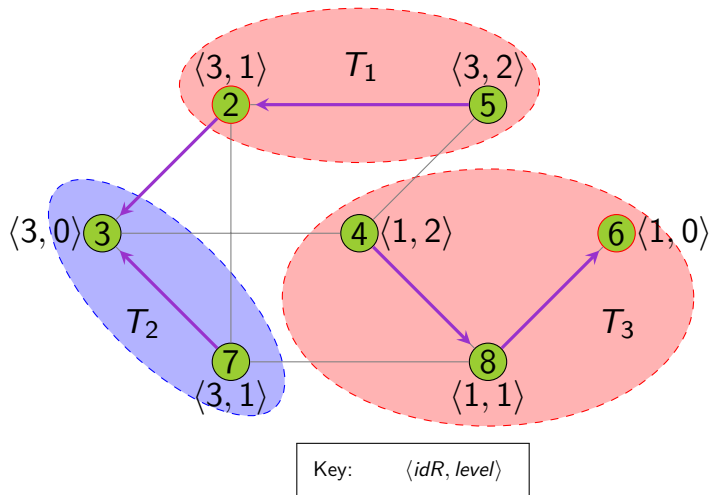




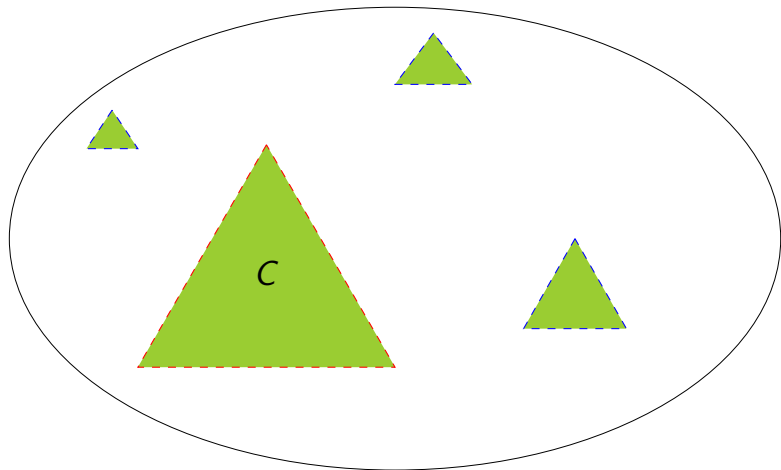
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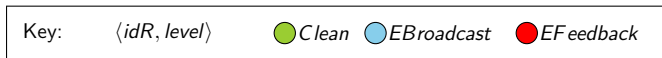
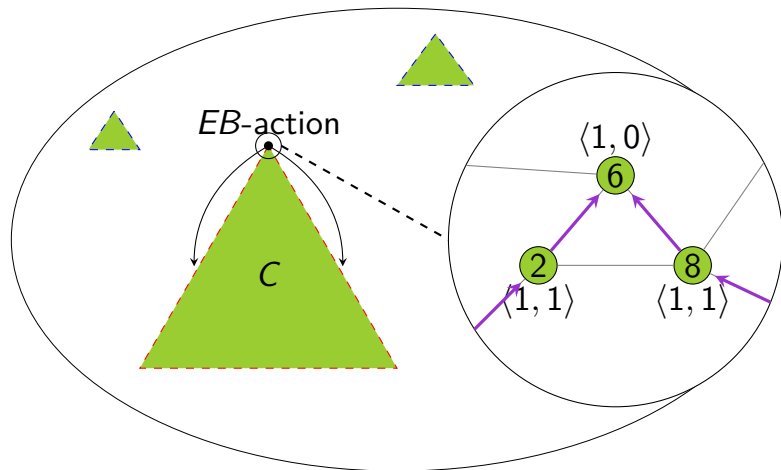


# Cleaning

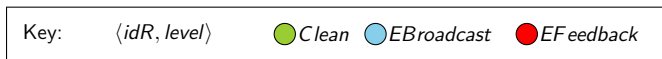
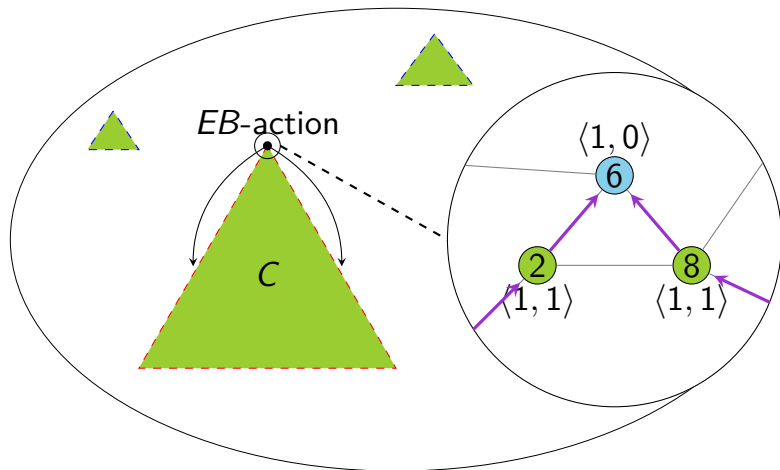


Key:  $\langle idR, level \rangle$     ● Clean    ● EBroadcast    ● EFeedback

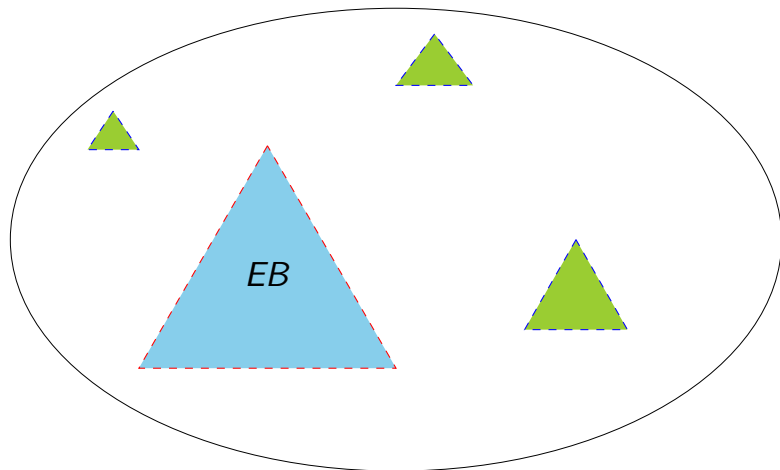
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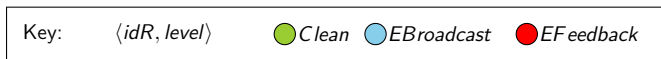
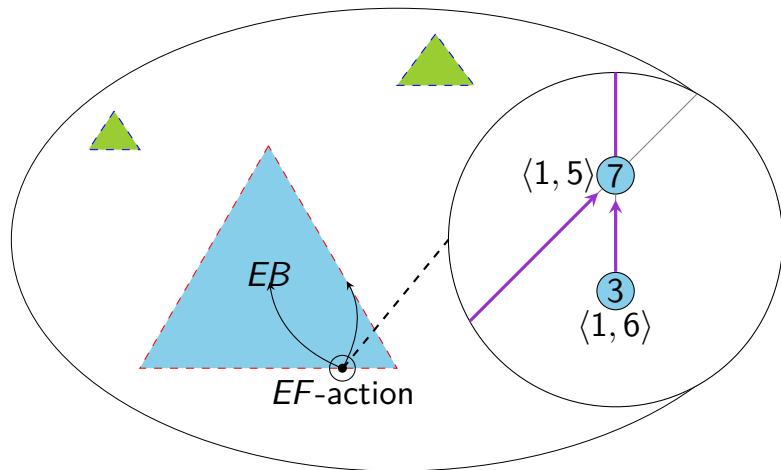


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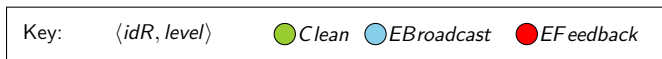
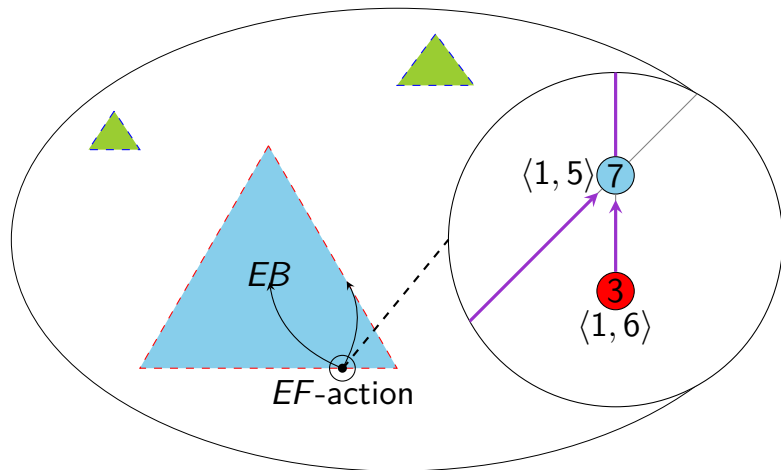


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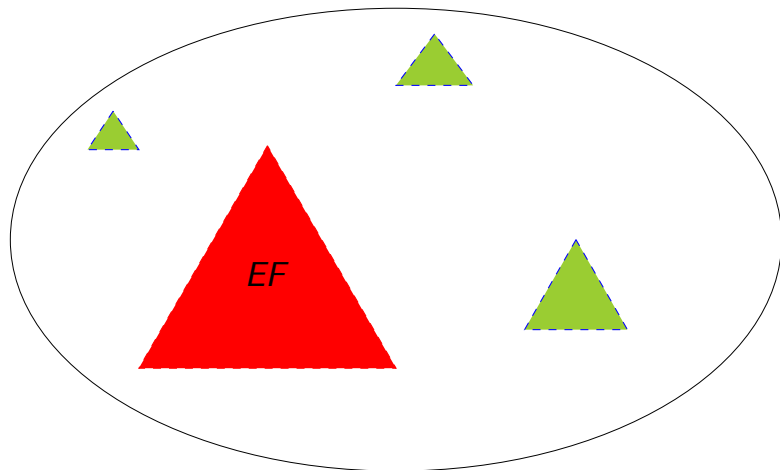


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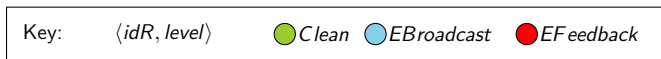
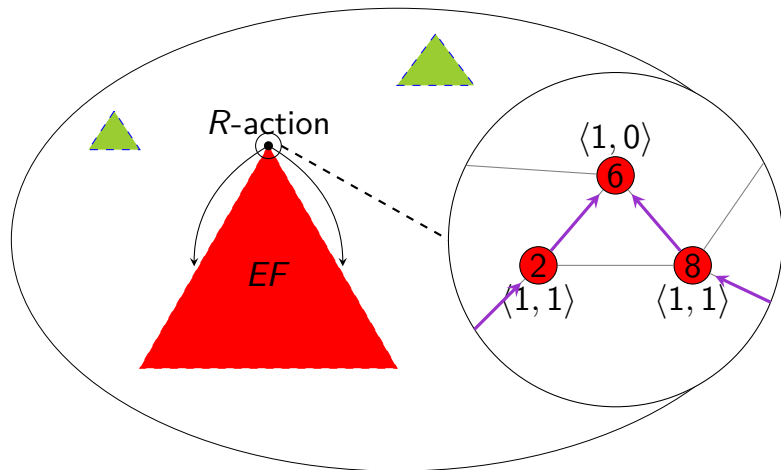


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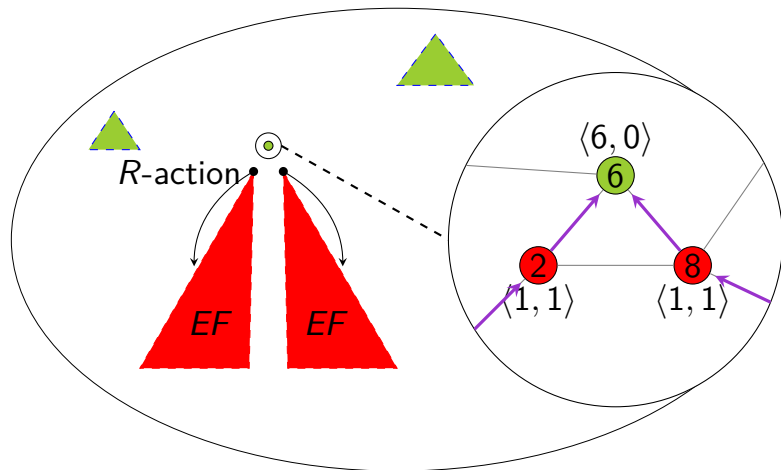


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- No alive abnormal tree created
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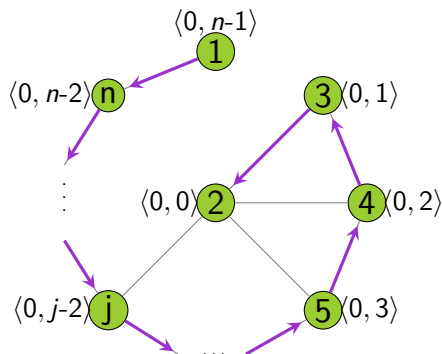
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$O(3n + \mathcal{D})$  rounds

# Lower Bound on the Worst Case Stabilization Time in Rounds

- $k$  links
- $j = k + 3$
- $\mathcal{D} = n - k$

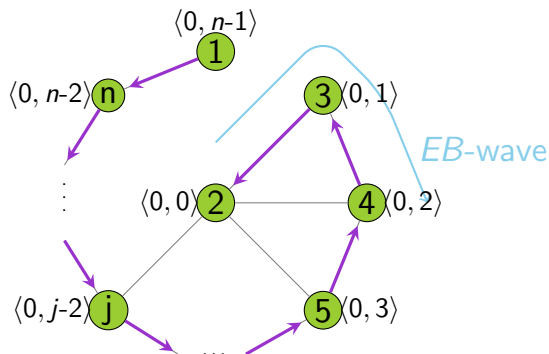


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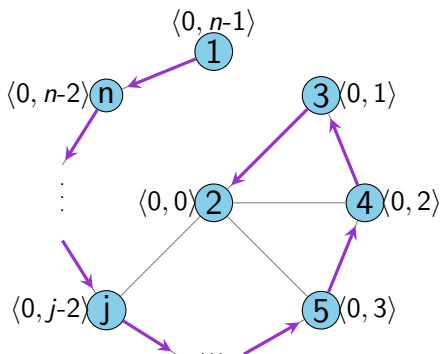
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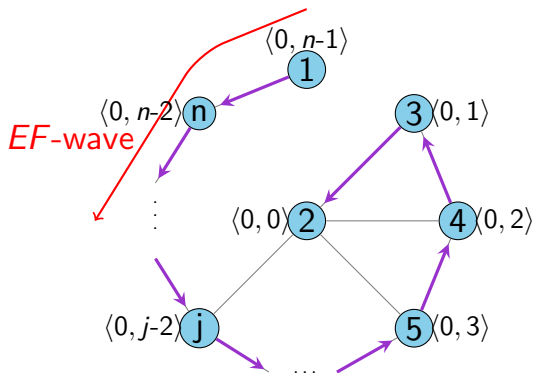


Key:  $\langle idR, level \rangle$     ● Clean    ● EBroadcast    ● EFeedback

$n$

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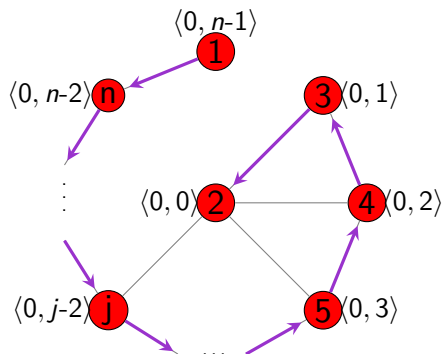


Key:	$\langle idR, level \rangle$	<span style="color: green;">●</span> Clean	<span style="color: lightblue;">●</span> EBroadcast	<span style="color: red;">●</span> EFeedback
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$n$

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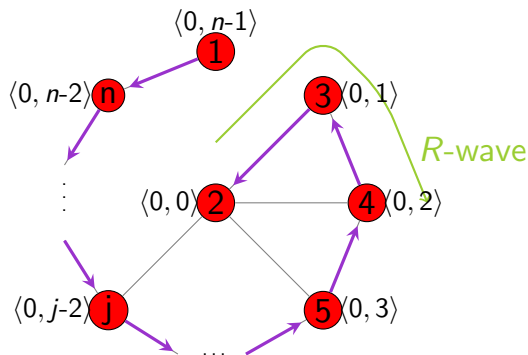


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$n + n$

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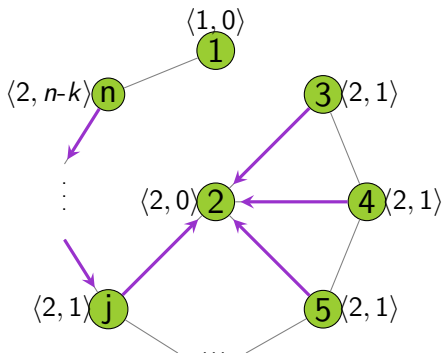


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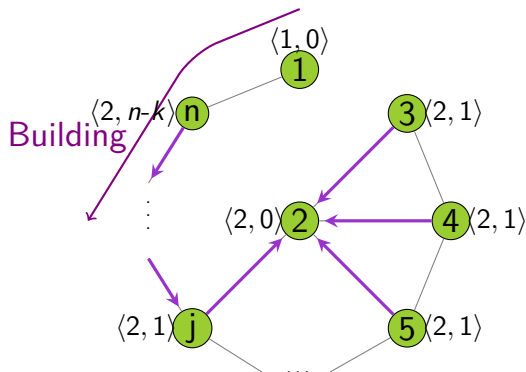


Key:  $\langle idR, level \rangle$     ● Clean    ● EBroadcast    ● EFeedback

$n + n + n$

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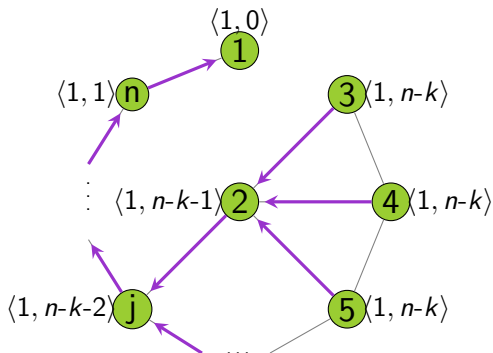


Key:  $\langle idR, level \rangle$     ● Clean    ● EBroadcast    ● EFeedback

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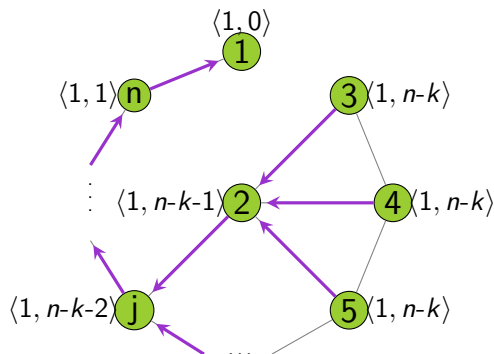
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$$n + n + n + (n - k)$$



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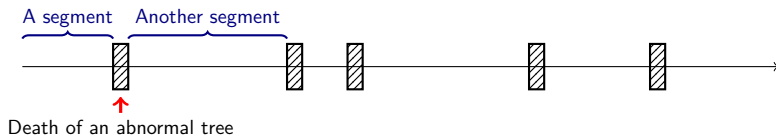
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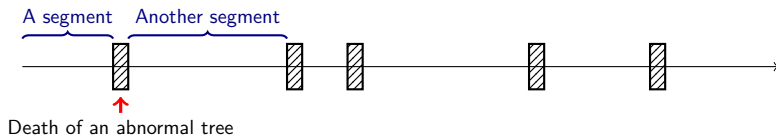
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$$\begin{aligned}
 & n + n + n + (n - k) \\
 = & \text{ exactly } 3n + \mathcal{D} \text{ rounds}
 \end{aligned}$$

# Stabilization Time in Steps

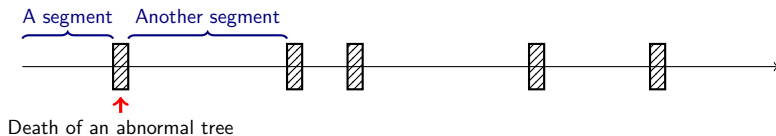


# Stabilization Time in Steps



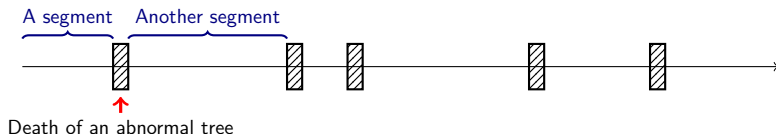
At most  $n$  alive abnormal trees + No alive abnormal tree created

# Stabilization Time in Steps



At most  $n$  alive abnormal trees      +      No alive abnormal tree created  
→ At most  $n + 1$  segments

# Stabilization Time in Steps

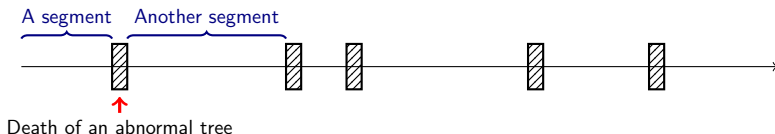


At most  $n$  alive abnormal trees + No alive abnormal tree created  
 $\rightarrow$  At most  $n + 1$  segments

## In a segment

$idR : 7 \xrightarrow{J\text{-action}} 5 \xrightarrow{J\text{-action}} 3 \xrightarrow{J\text{-action}} 2 \xrightarrow{EB\text{-action}} \xrightarrow{EF\text{-action}} \xrightarrow{R\text{-action}} 7 \xrightarrow{J\text{-action}} 3$   
 Death of an abnormal tree = End of the segment

# Stabilization Time in Steps



At most  $n$  alive abnormal trees + No alive abnormal tree created  
 $\rightarrow$  At most  $n + 1$  segments

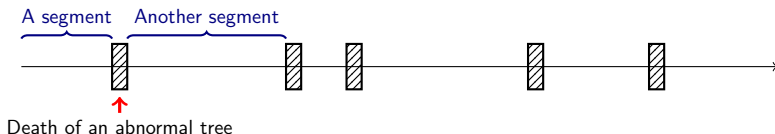
## In a segment

$idR : 7 \xrightarrow{J\text{-action}} 5 \xrightarrow{J\text{-action}} 3 \xrightarrow{J\text{-action}} 2 \xrightarrow{EB\text{-action}} \xrightarrow{EF\text{-action}} \xrightarrow{R\text{-action}} 7 \xrightarrow{J\text{-action}} 3$

Death of an abnormal tree = End of the segment

- $n - 1$   $J$ -actions
- 1  $EB$ -action
- 1  $EF$ -action
- 1  $R$ -action

# Stabilization Time in Steps



At most  $n$  alive abnormal trees + No alive abnormal tree created  
 $\longrightarrow$  At most  $n + 1$  segments

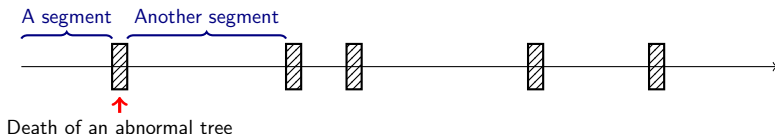
## In a segment

$idR : 7 \xrightarrow{J\text{-action}} 5 \xrightarrow{J\text{-action}} 3 \xrightarrow{J\text{-action}} 2 \xrightarrow{EB\text{-action}} \xrightarrow{EF\text{-action}} \xrightarrow{R\text{-action}} 7 \xrightarrow{J\text{-action}} 3$

Death of an abnormal tree = End of the segment

- $n - 1$   $J$ -actions
  - 1  $EB$ -action
  - 1  $EF$ -action
  - 1  $R$ -action
- $\Rightarrow O(n)$  actions per process

# Stabilization Time in Steps



At most  $n$  alive abnormal trees + No alive abnormal tree created  
 $\rightarrow$  At most  $n + 1$  segments

## In a segment

$idR : 7 \xrightarrow{J\text{-action}} 5 \xrightarrow{J\text{-action}} 3 \xrightarrow{J\text{-action}} 2 \xrightarrow{EB\text{-action}} \xrightarrow{EF\text{-action}} \xrightarrow{R\text{-action}} 7 \xrightarrow{J\text{-action}} 3$

Death of an abnormal tree = End of the segment

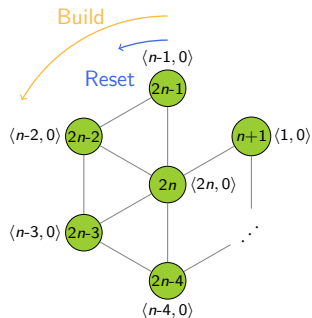
- $n - 1$   $J$ -actions
  - 1  $EB$ -action
  - 1  $EF$ -action
  - 1  $R$ -action
- $\Rightarrow O(n)$  actions per process

$O(n^3)$  steps

Lower Bound:  $\frac{n^3}{6} + \frac{5}{2}n^2 - \frac{11}{3}n + 2$  steps    Upper Bound:  $\frac{n^3}{2} + 2n^2 + \frac{n}{2} + 1$  steps



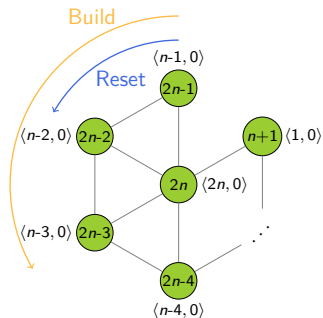
# Lower Bound on the Worst Case Stabilization Time in Steps



Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

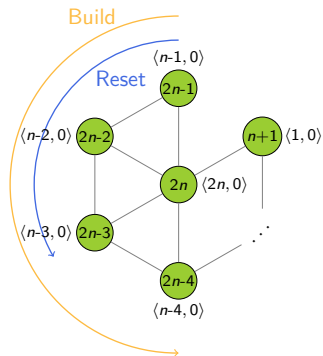
# Lower Bound on the Worst Case Stabilization Time in Steps



Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

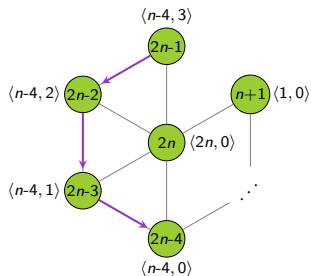
# Lower Bound on the Worst Case Stabilization Time in Steps



Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

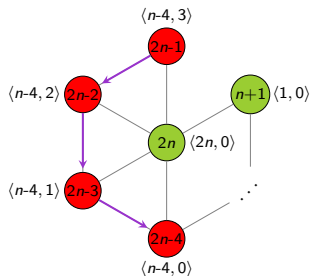
# Lower Bound on the Worst Case Stabilization Time in Steps



Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

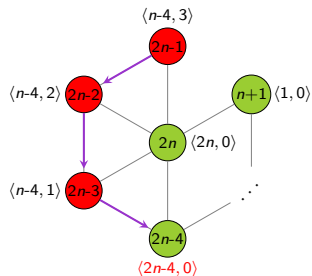
# Lower Bound on the Worst Case Stabilization Time in Steps



Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

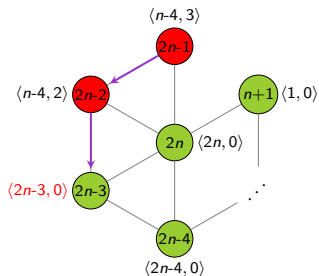
# Lower Bound on the Worst Case Stabilization Time in Steps



Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

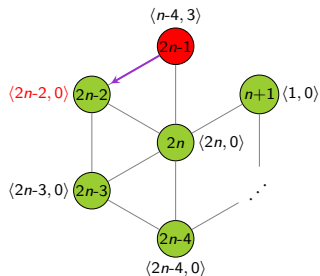
# Lower Bound on the Worst Case Stabilization Time in Steps



Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

# Lower Bound on the Worst Case Stabilization Time in Steps



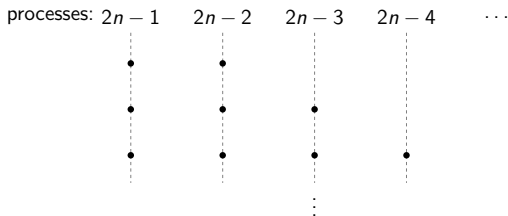
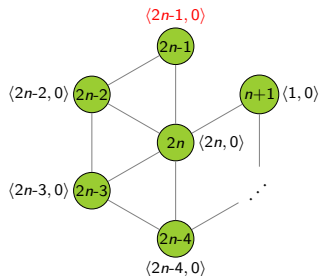
Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback



# Lower Bound on the Worst Case Stabilization Time in Steps

## Case of the reset of $2n - 4$

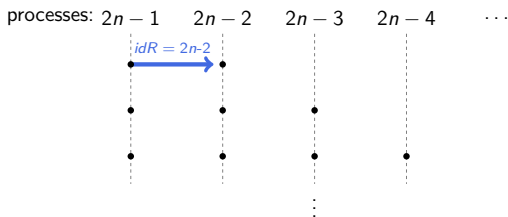
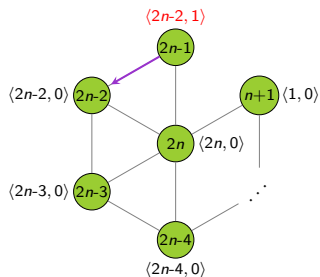


Key:  $\langle idR, level \rangle$

● *Clean* ● *EBroadcast* ● *EFeedback*

# Lower Bound on the Worst Case Stabilization Time in Steps

## Case of the reset of $2n - 4$

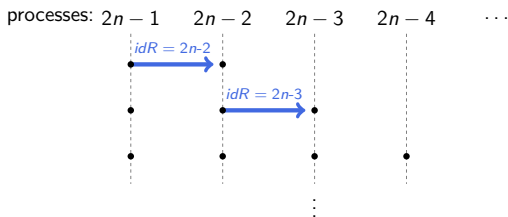
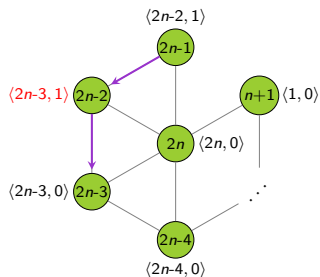


Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

# Lower Bound on the Worst Case Stabilization Time in Steps

## Case of the reset of $2n - 4$

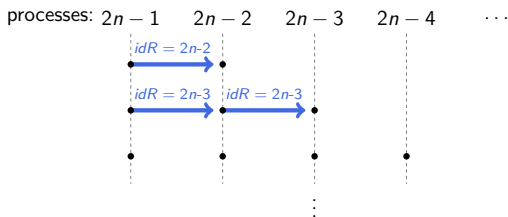
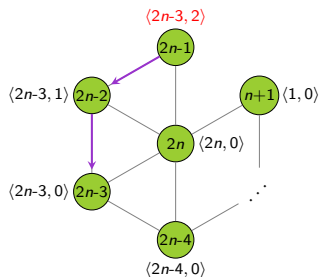


Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

# Lower Bound on the Worst Case Stabilization Time in Steps

## Case of the reset of $2n - 4$

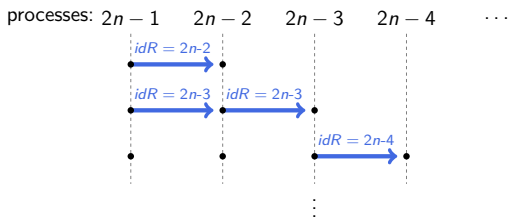
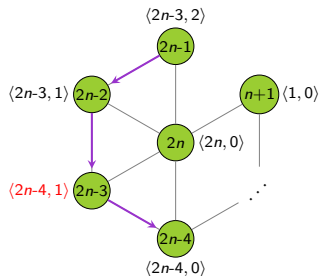


Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

# Lower Bound on the Worst Case Stabilization Time in Steps

## Case of the reset of $2n - 4$

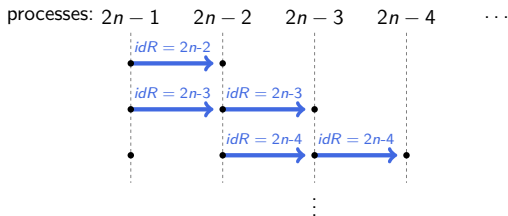
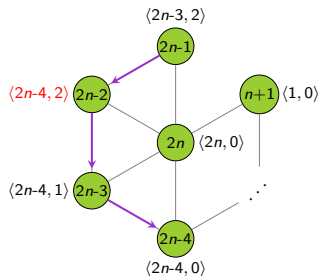


Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

# Lower Bound on the Worst Case Stabilization Time in Steps

## Case of the reset of $2n - 4$

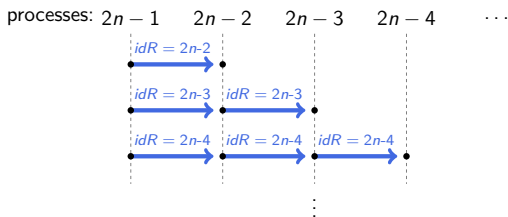
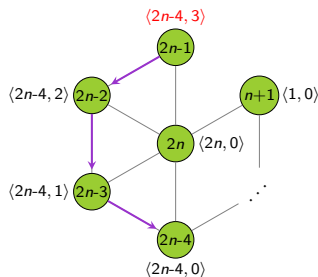


Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

# Lower Bound on the Worst Case Stabilization Time in Steps

## Case of the reset of $2n - 4$

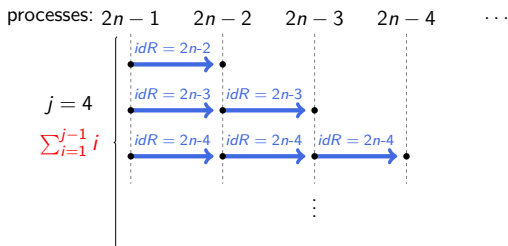
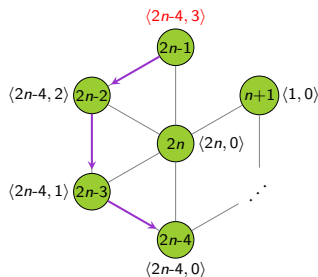


Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

# Lower Bound on the Worst Case Stabilization Time in Steps

## Case of the reset of $2n - 4$



Key:  $\langle idR, level \rangle$

● Clean ● EBroadcast ● EFeedback

$$\Theta(n) \text{ reset} \Rightarrow \sum_{j=1}^n \sum_{i=1}^{j-1} i \Rightarrow \Theta(n^3) \text{ steps}$$



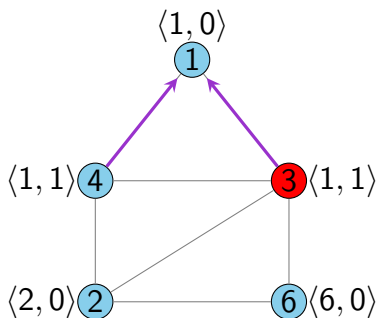
# Analytical Study of Datta *et al*, 2011<sup>3</sup>

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<sup>3</sup>Datta, Larmore, and Vemula. Self-stabilizing Leader Election in Optimal Space under an Arbitrary Scheduler. 2011

# Principles

## Join a tree



Key:

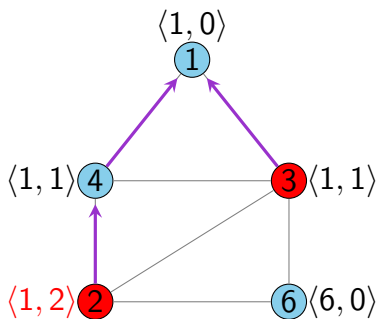
$\langle idR, level \rangle$

● Can be joined

● Cannot be joined

# Principles

## Join a tree



Key:

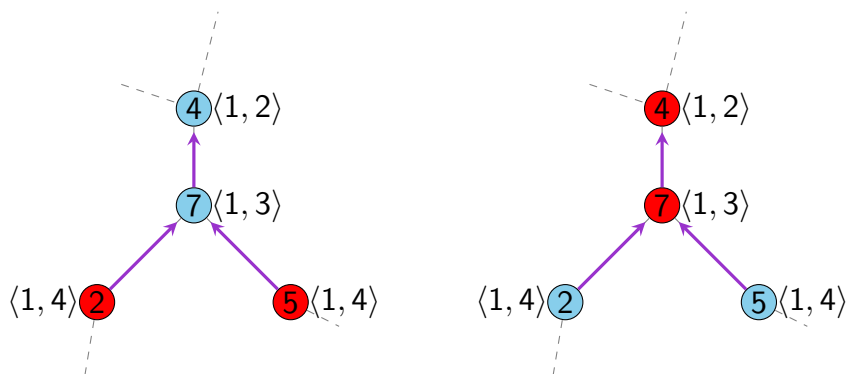
$\langle idR, level \rangle$

● Can be joined

● Cannot be joined

# Principles

## Change of color



Key:

$\langle idR, level \rangle$



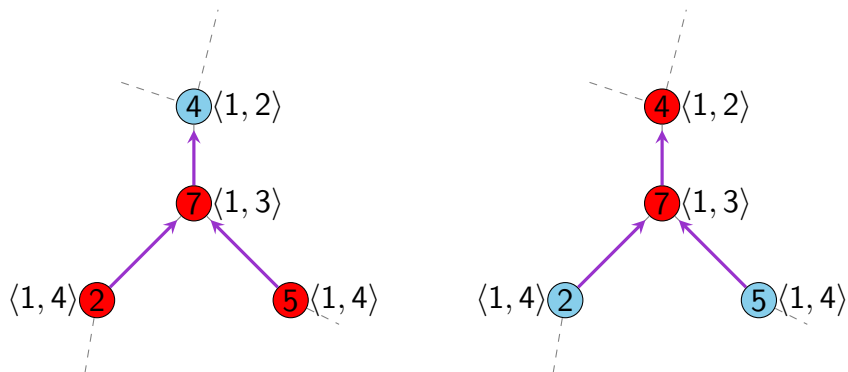
Can be joined



Cannot be joined

# Principles

## Change of color



Key:

$\langle idR, level \rangle$



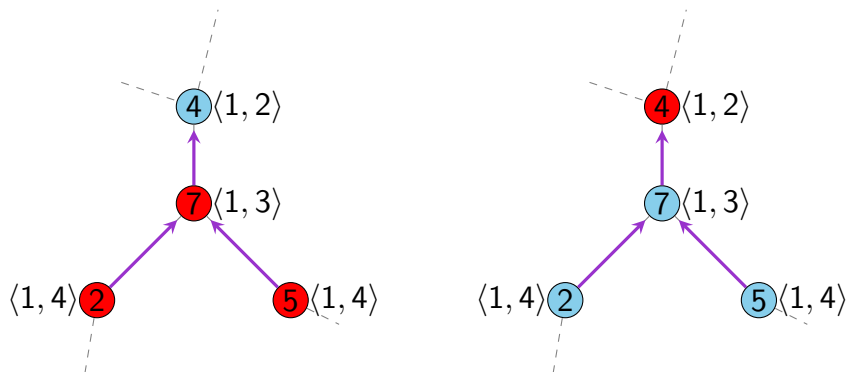
Can be joined



Cannot be joined

# Principles

## Change of color



Key:

$\langle idR, level \rangle$



Can be joined

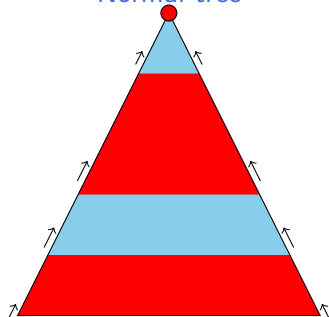


Cannot be joined

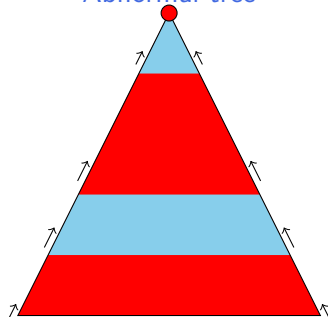
# Principles

## Color Waves Absorption

Normal tree



Abnormal tree



Key:

$\langle idR, level \rangle$

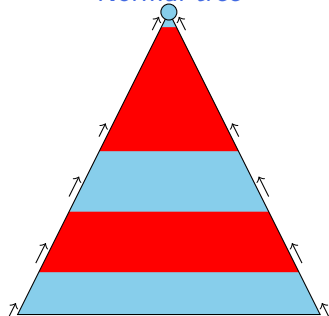
● Can be joined

● Cannot be joined

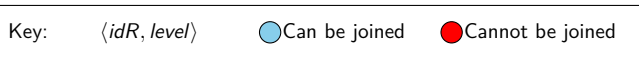
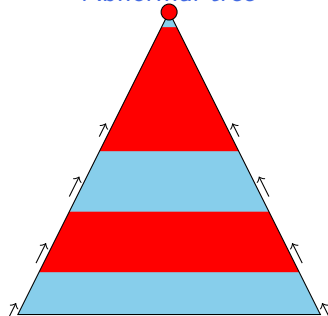
# Principles

## Color Waves Absorption

Normal tree



Abnormal tree

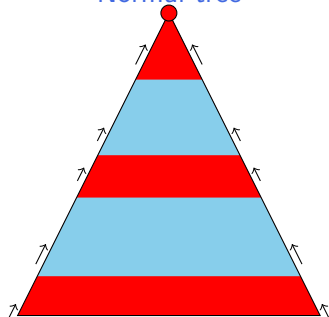




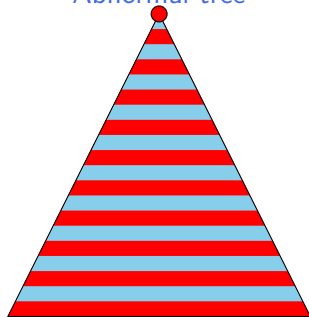
# Principles

## Color Waves Absorption

Normal tree



Abnormal tree



Key:

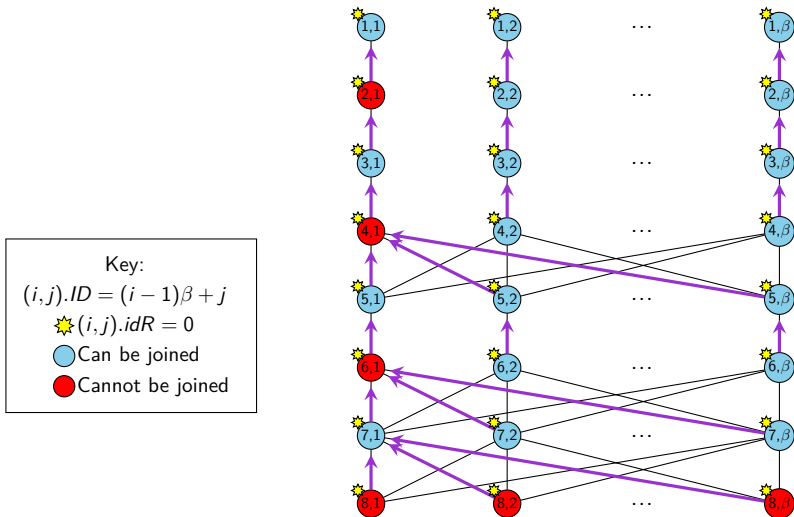
$\langle idR, level \rangle$

● Can be joined

● Cannot be joined

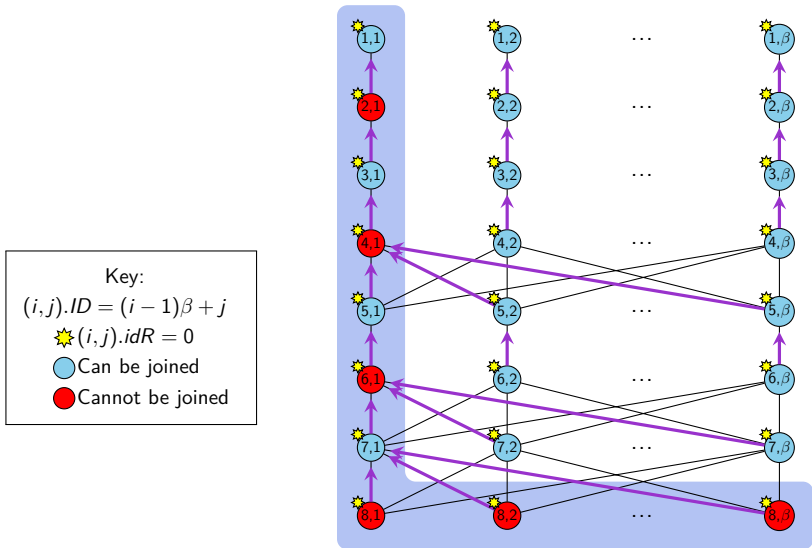
# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$



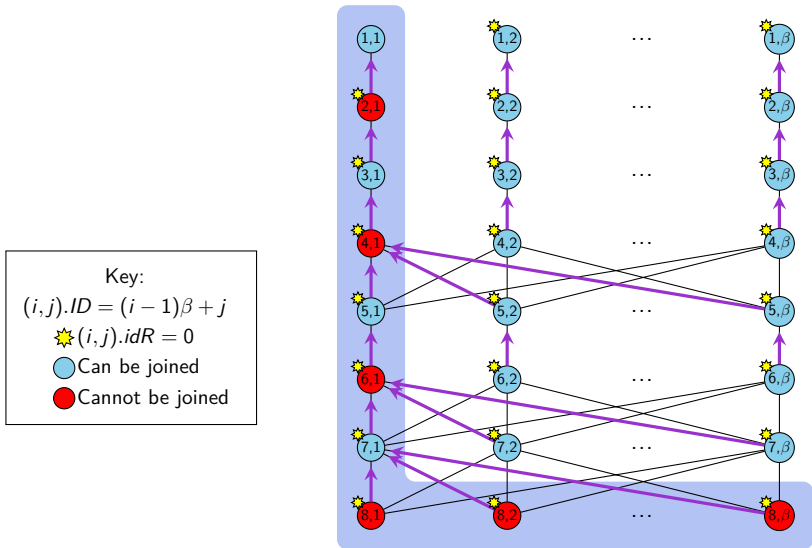
# Datta et al, 2011

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# Datta et al, 2011

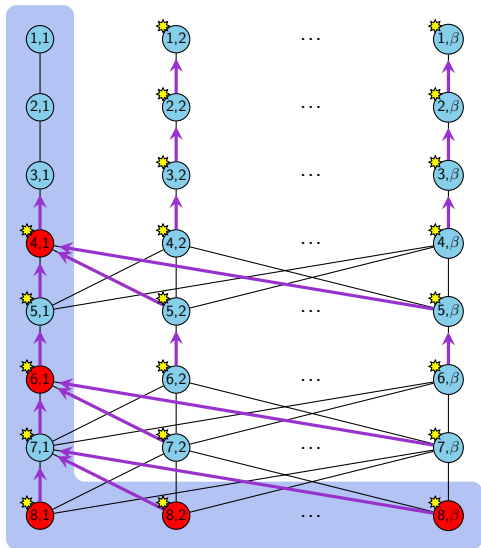
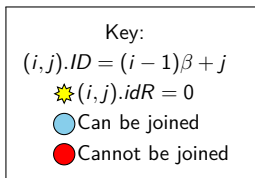
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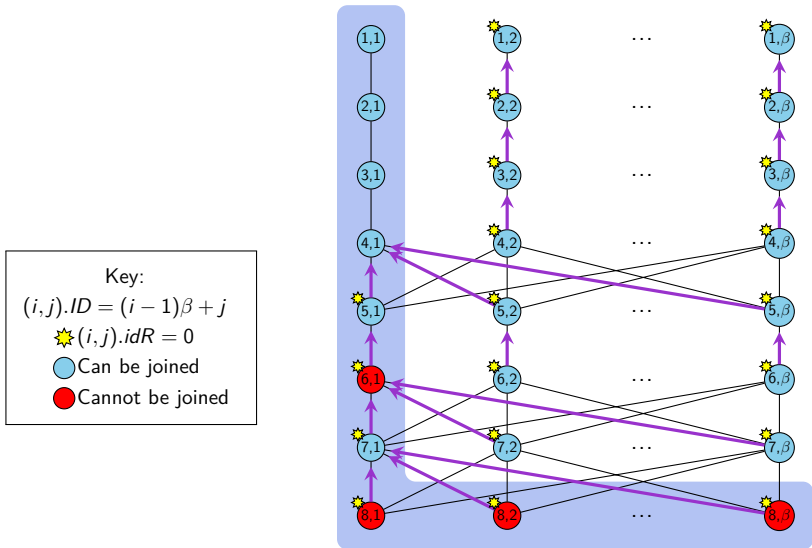
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Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$



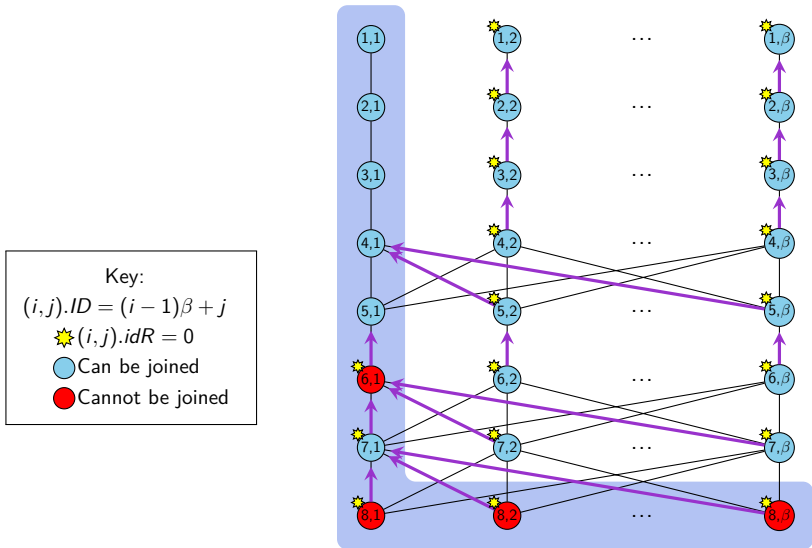
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# Datta et al, 2011

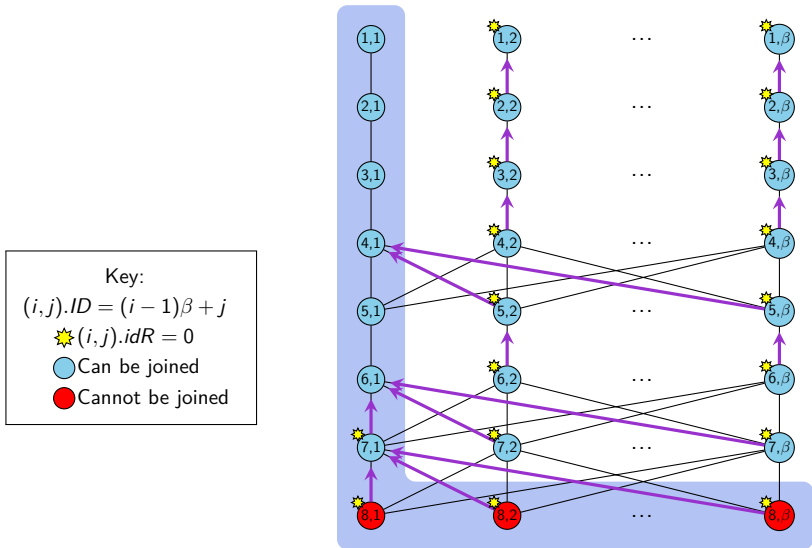
Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$





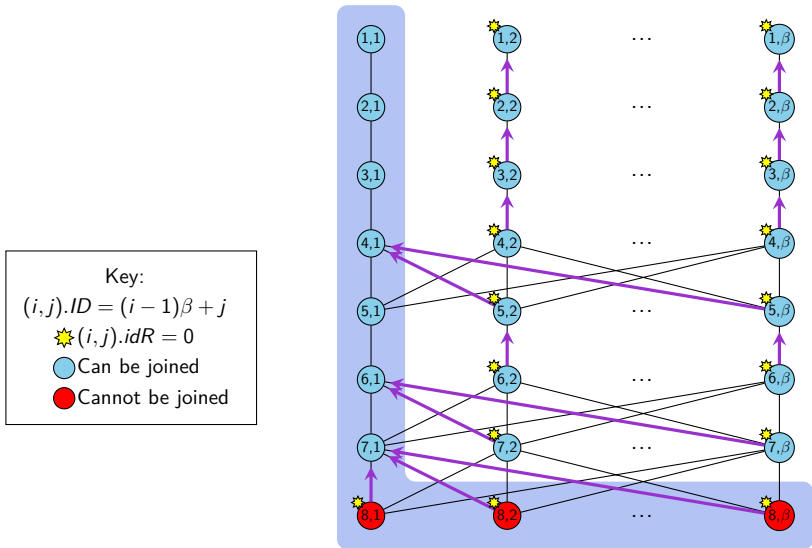
# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$



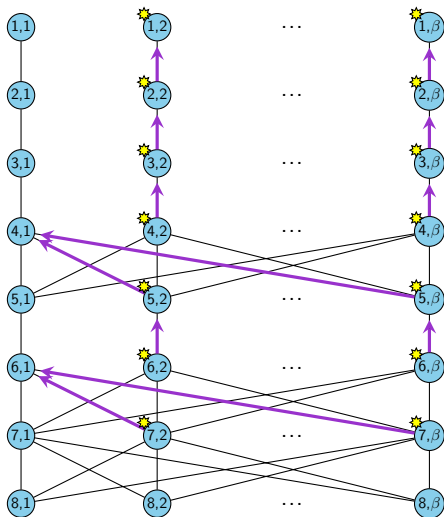
# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta$

Key:

- $(i,j).ID = (i-1)\beta + j$
- ★  $(i,j).idR = 0$
- Can be joined
- Cannot be joined

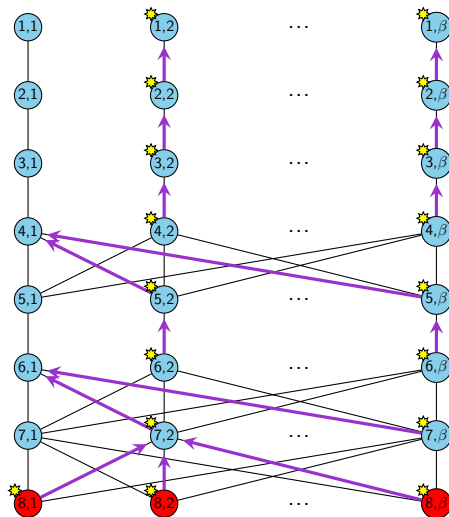
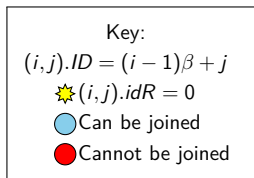




# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta$



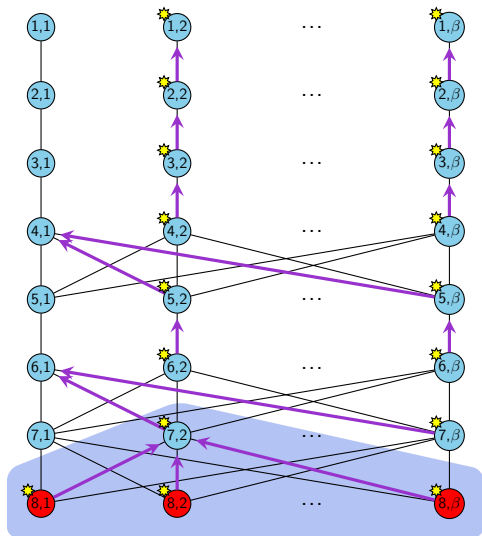
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Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta$

Key:

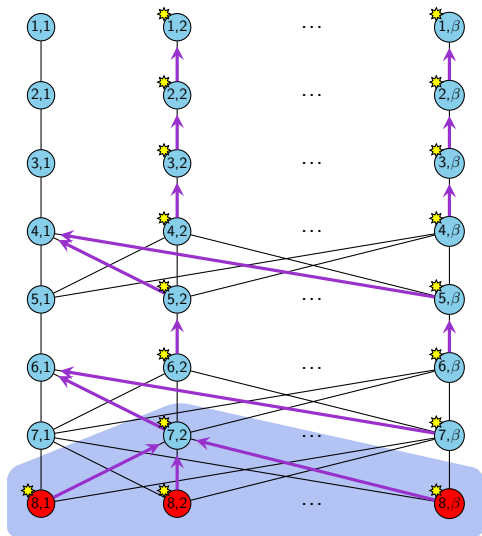
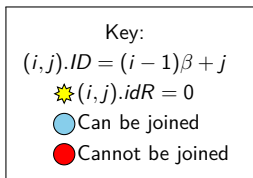
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# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

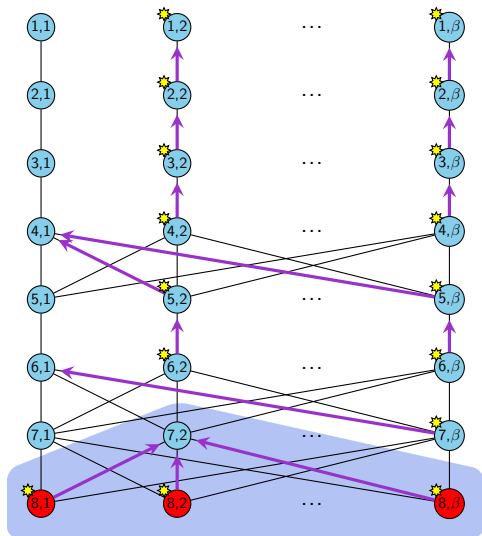
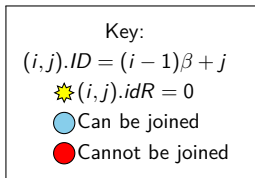
$\beta$



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta$

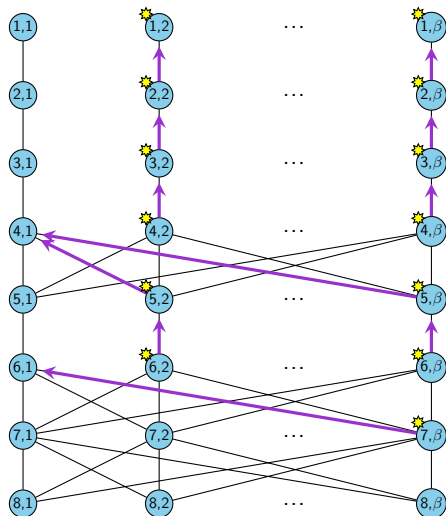
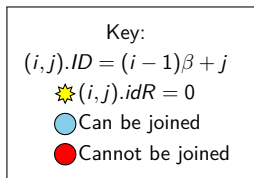




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Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta$

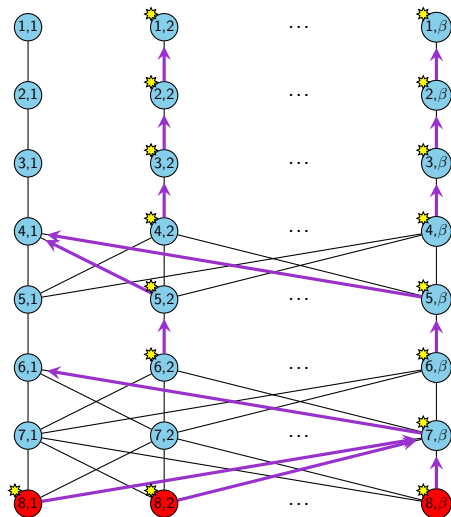
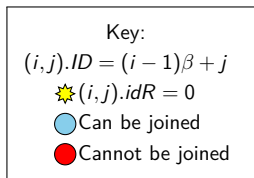




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$\beta$



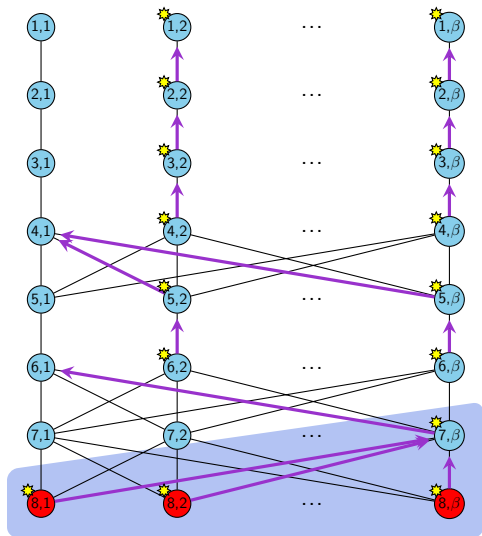
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Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta$

Key:

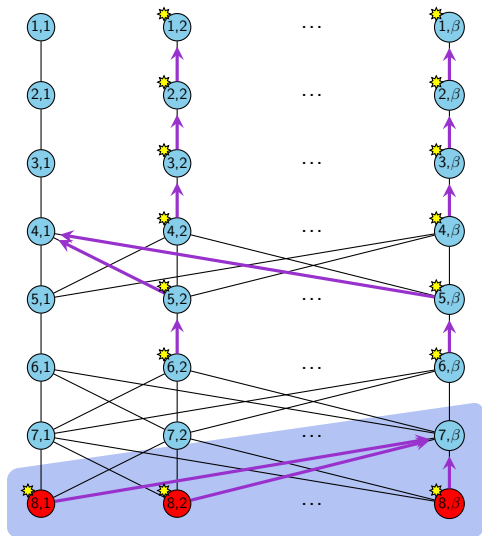
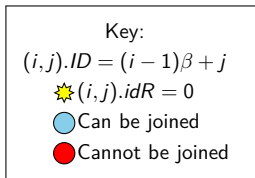
- $(i,j).ID = (i-1)\beta + j$
- ★  $(i,j).idR = 0$
- Can be joined
- Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta$





# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

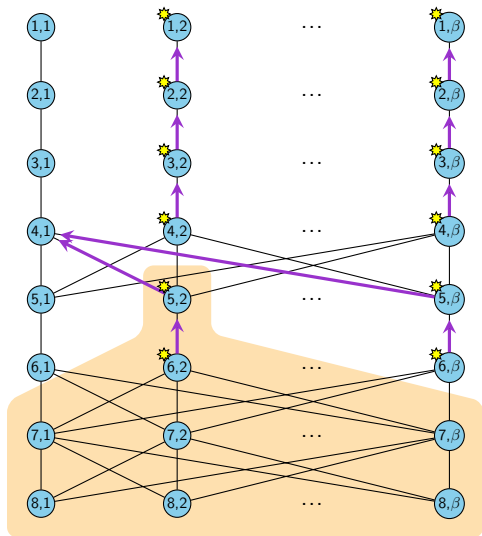
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

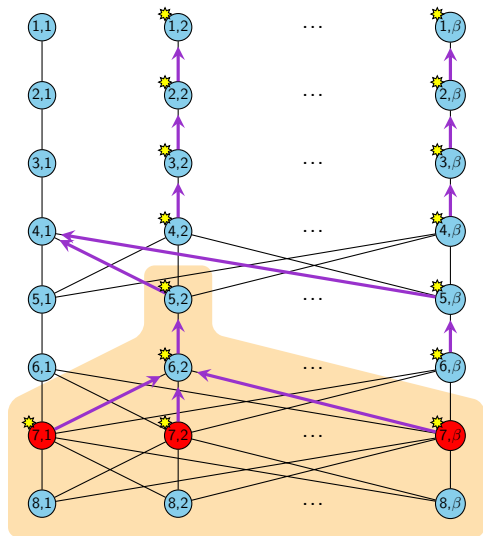
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined





# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

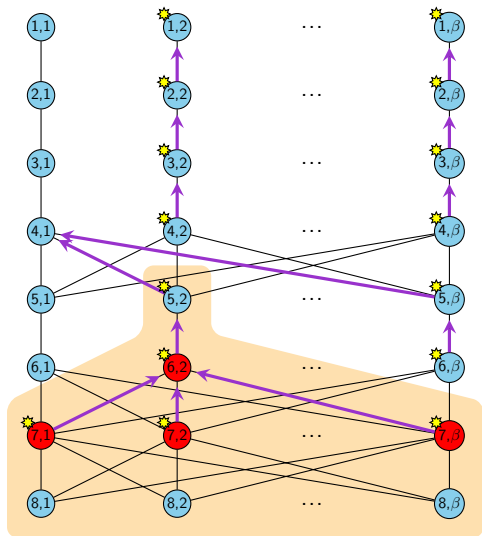
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

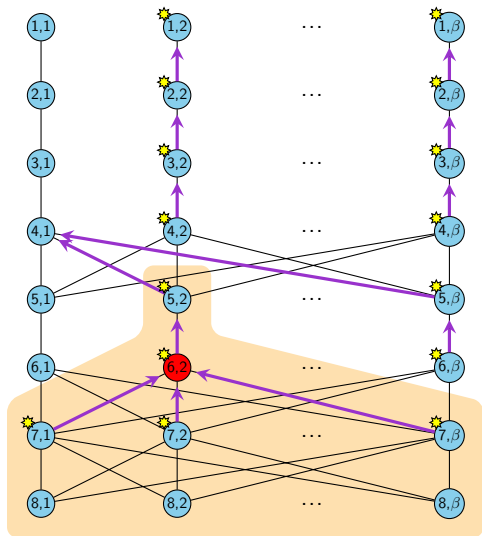
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

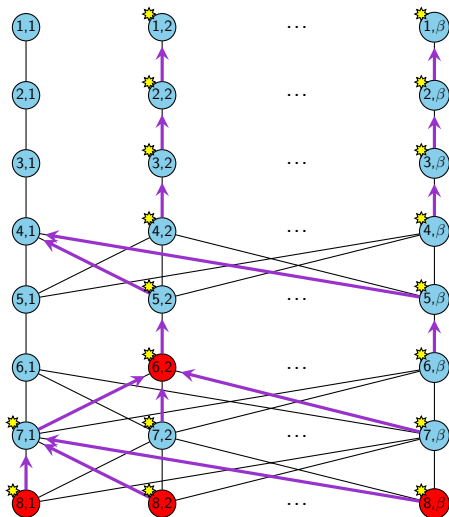
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

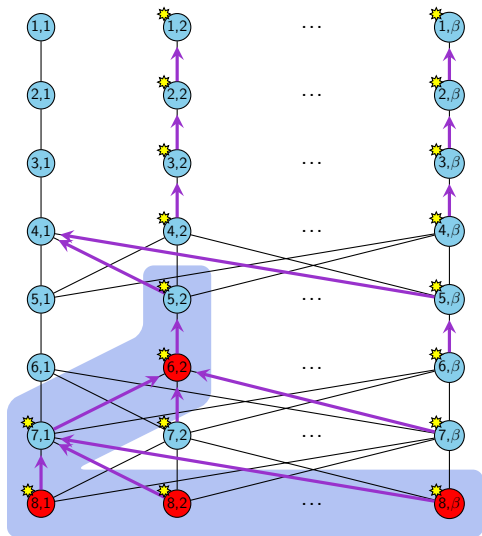
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

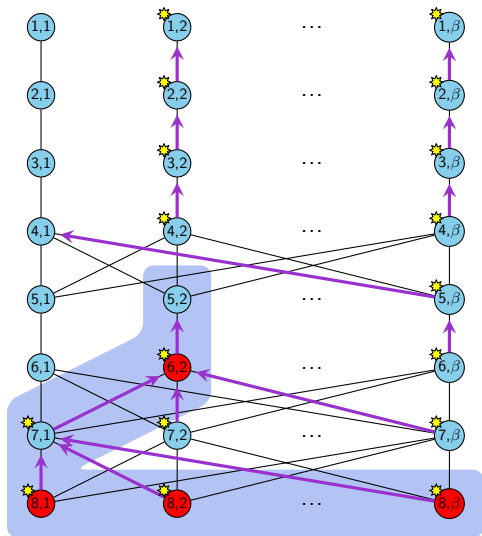
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

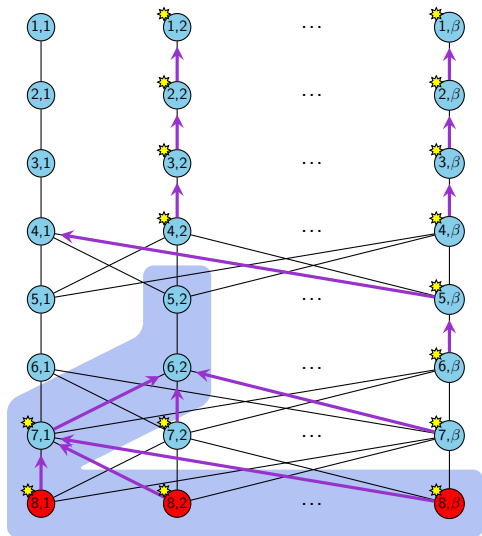
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined





# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

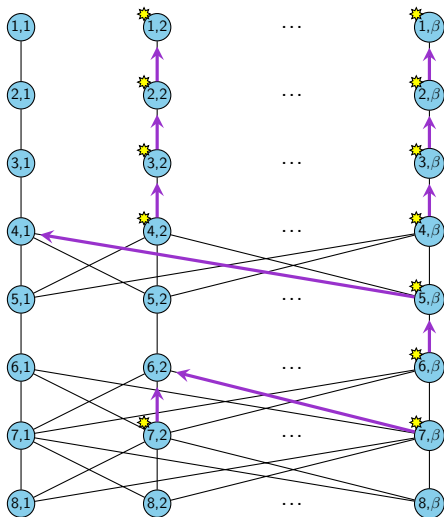
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined





# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

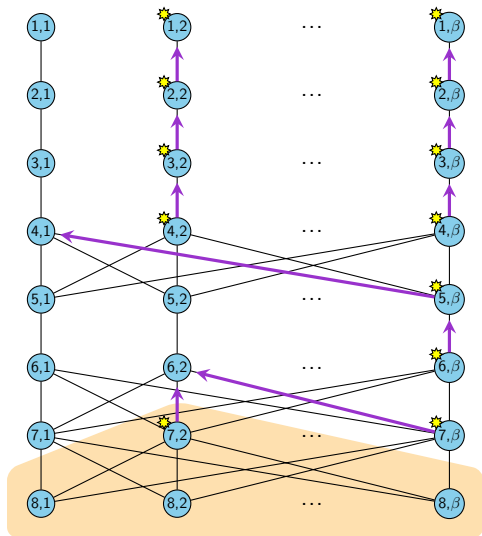
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

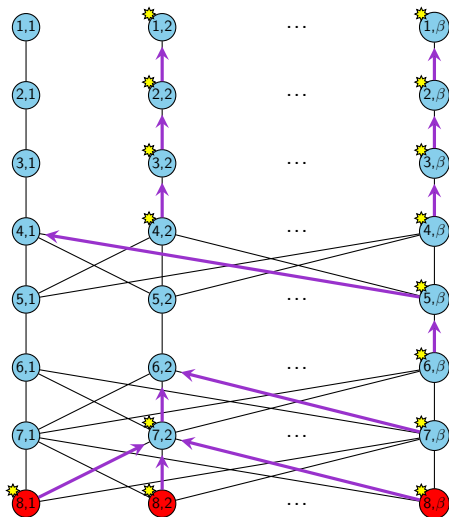
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



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Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

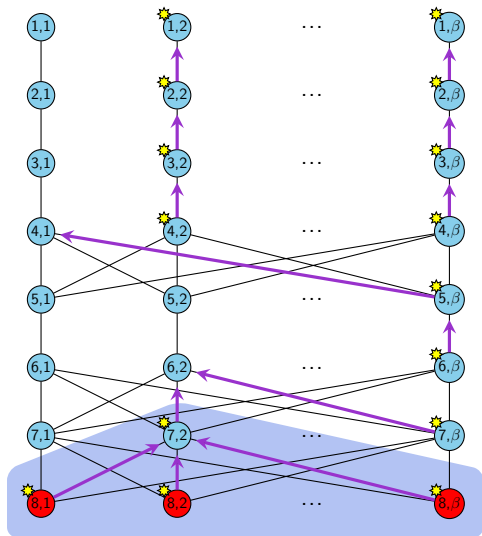
Key:

$(i, j).ID = (i - 1)\beta + j$

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Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

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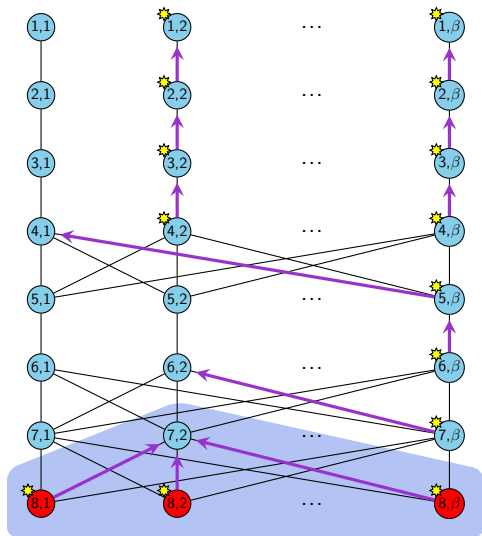
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

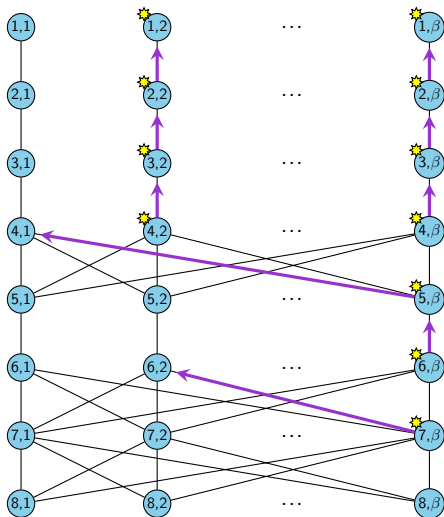
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

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# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

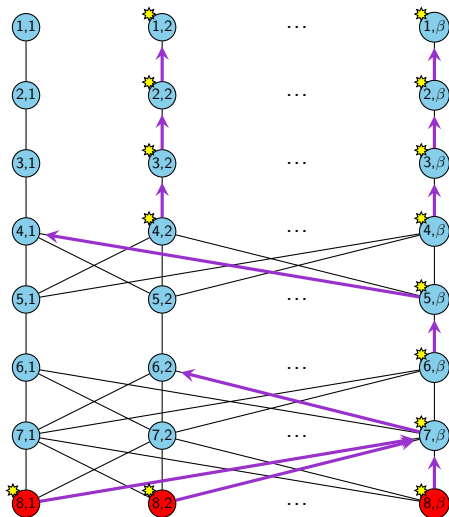
Key:

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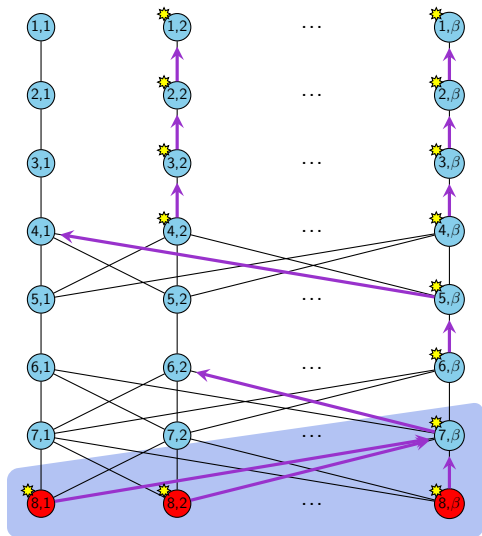
Key:

$(i, j).ID = (i - 1)\beta + j$

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● Can be joined

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# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

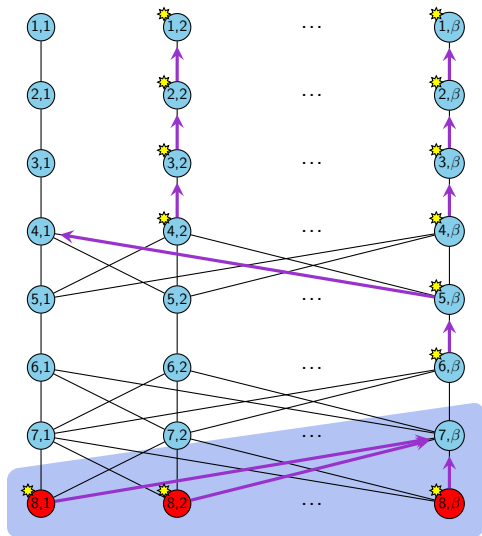
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

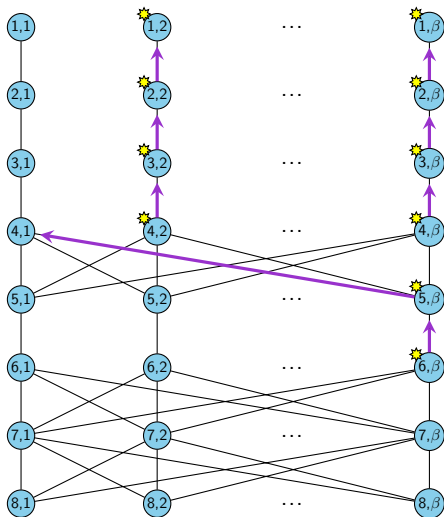
Key:

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★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

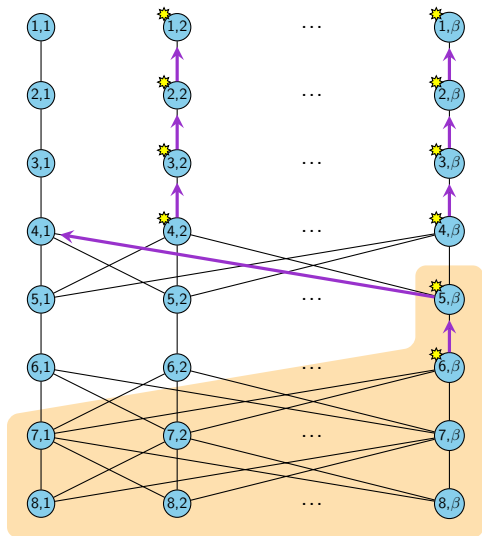
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

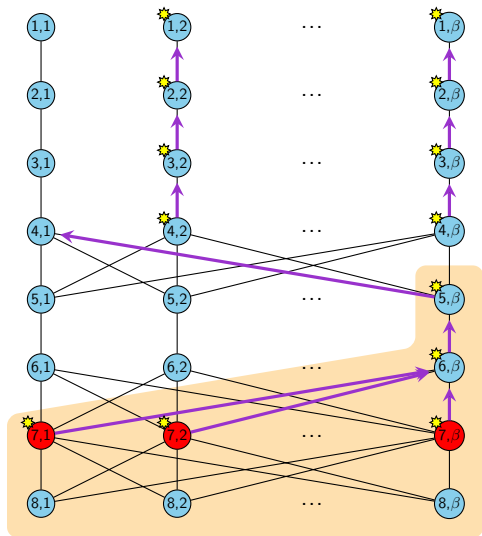
Key:

$(i, j).ID = (i - 1)\beta + j$

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# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

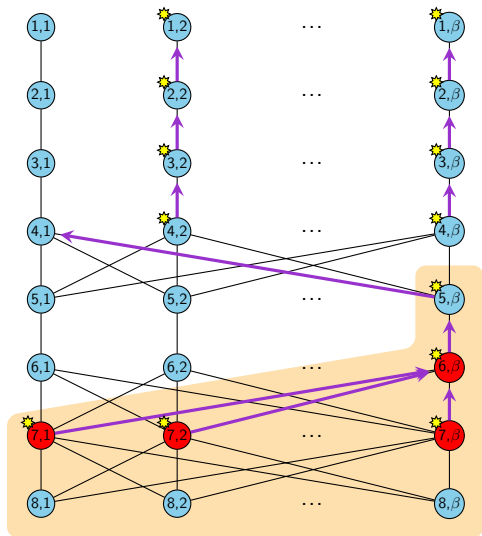
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined

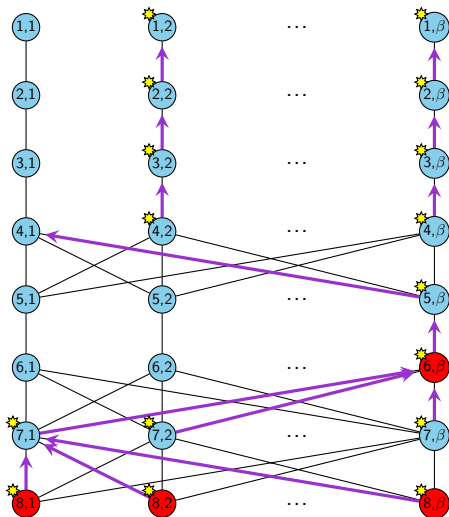
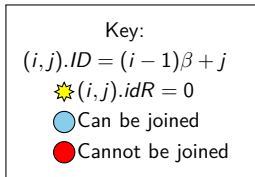




# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

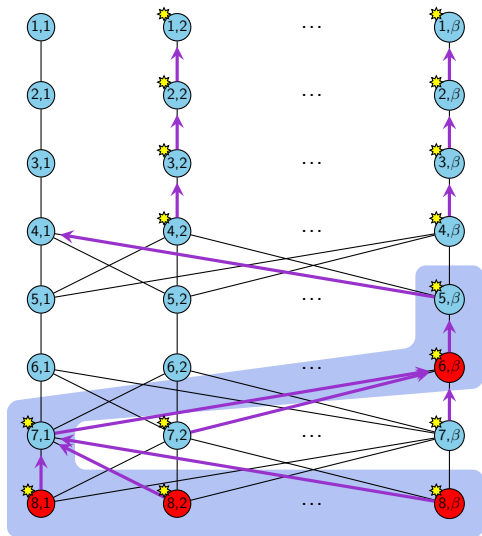
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined







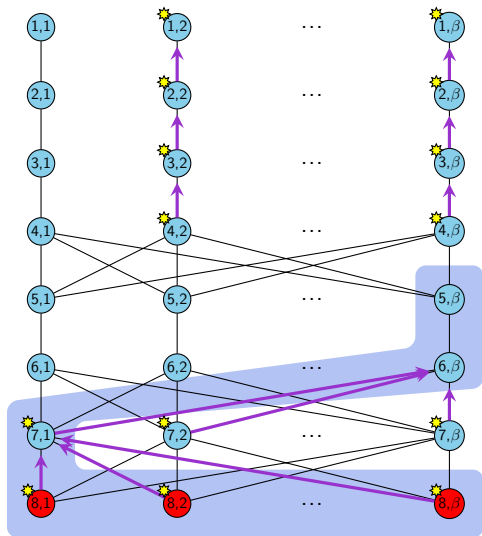
# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

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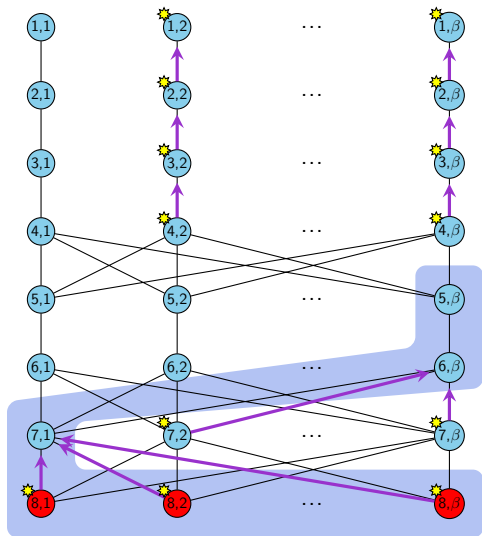
# Datta et al, 2011

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# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

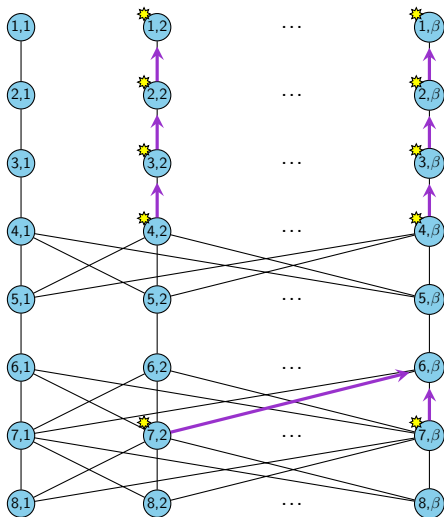
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



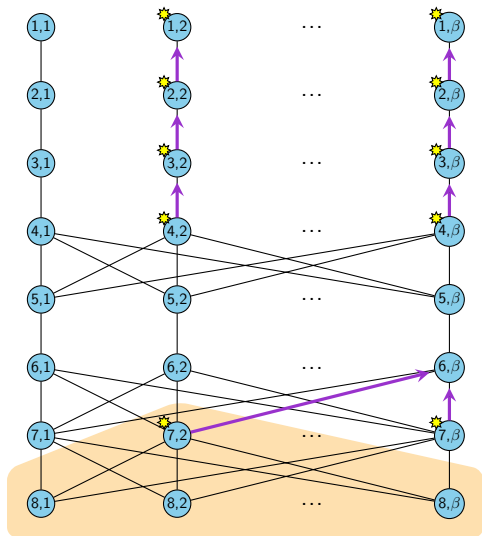
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# Datta et al, 2011

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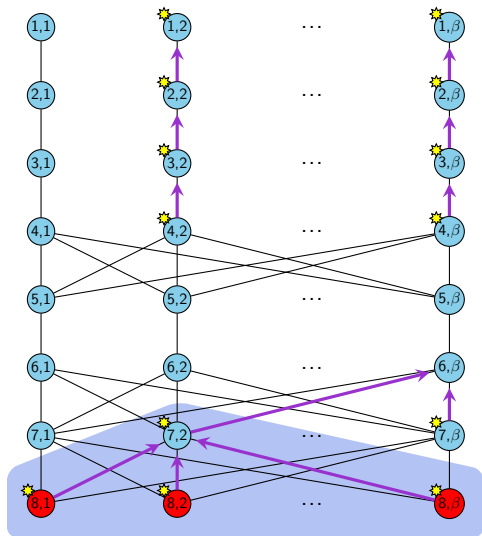
Key:

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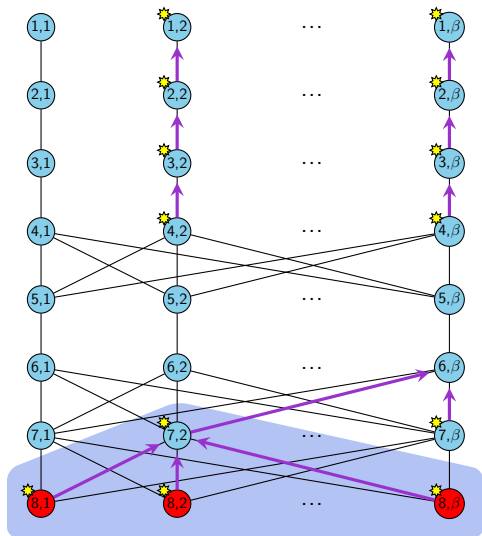
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# Datta et al, 2011

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$\beta^2$

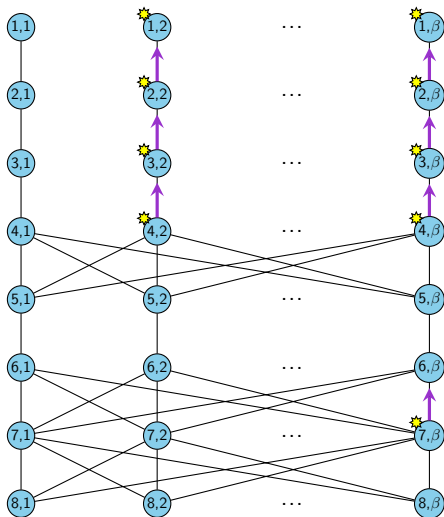
Key:

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● Can be joined

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# Datta et al, 2011

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$\beta^2$

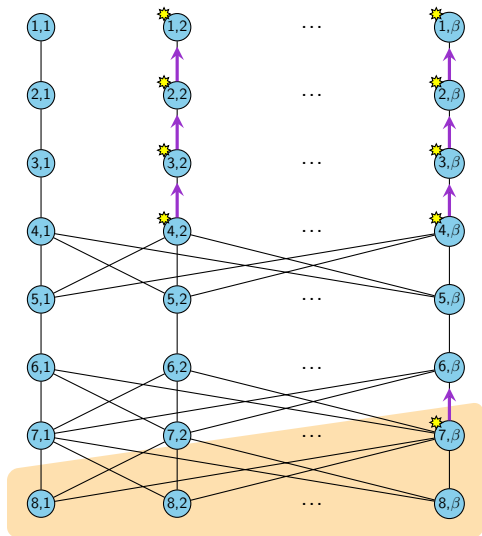
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

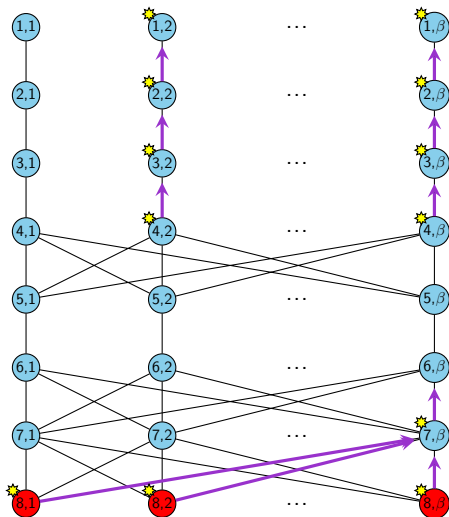
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

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# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^2$

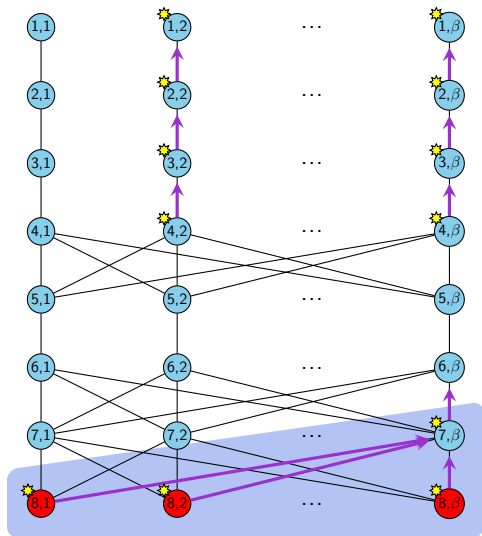
Key:

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● Can be joined

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# Datta et al, 2011

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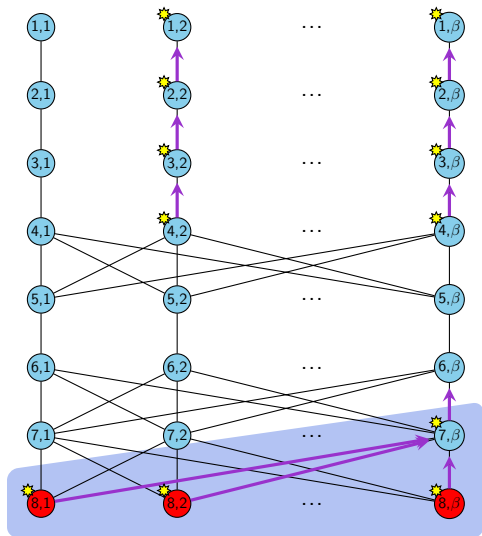
Key:

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● Can be joined

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# Datta et al, 2011

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$\beta^2$

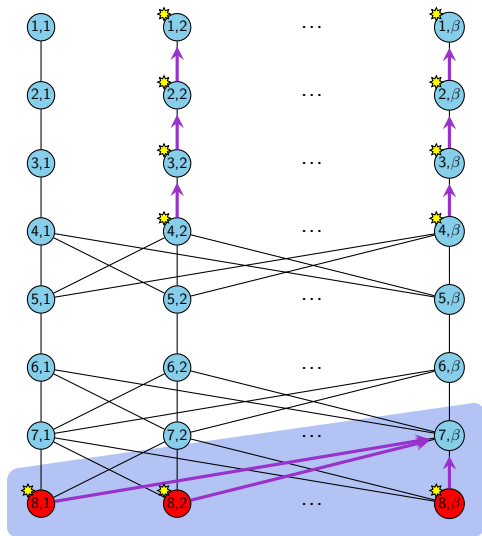
Key:

$(i, j).ID = (i - 1)\beta + j$

★  $(i, j).idR = 0$

● Can be joined

● Cannot be joined



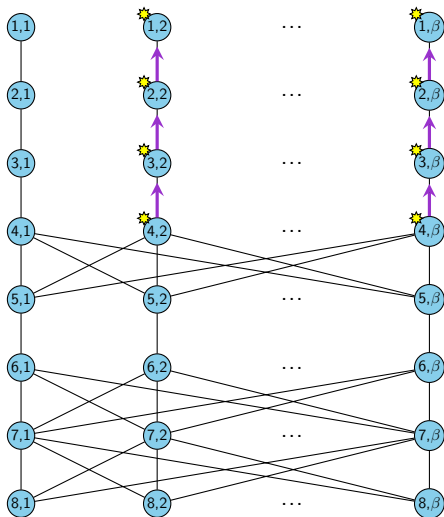
# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^3$

Key:

- $(i,j).ID = (i-1)\beta + j$
- ★  $(i,j).idR = 0$
- Can be joined
- Cannot be joined







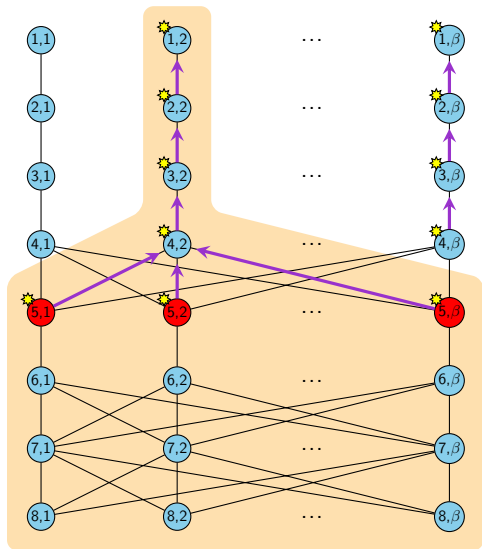
# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^3$

Key:

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- ★  $(i, j).idR = 0$
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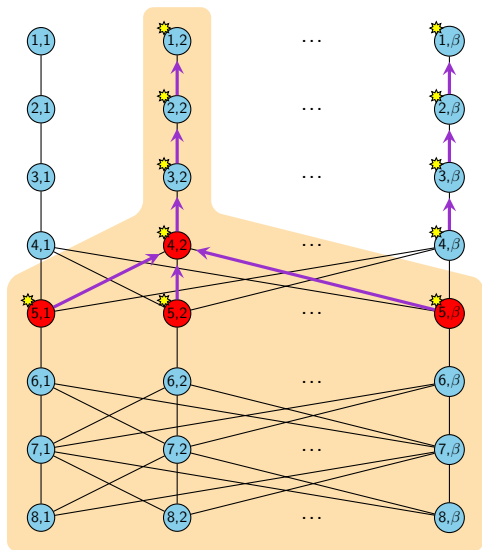
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Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^3$

Key:

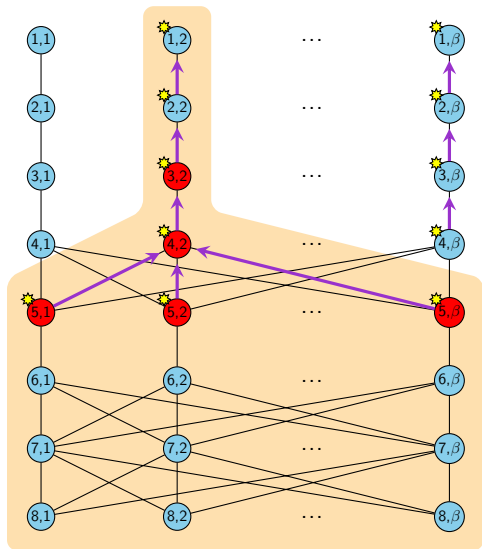
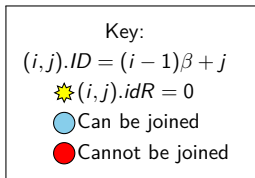
- $(i, j).ID = (i - 1)\beta + j$
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$\beta^3$



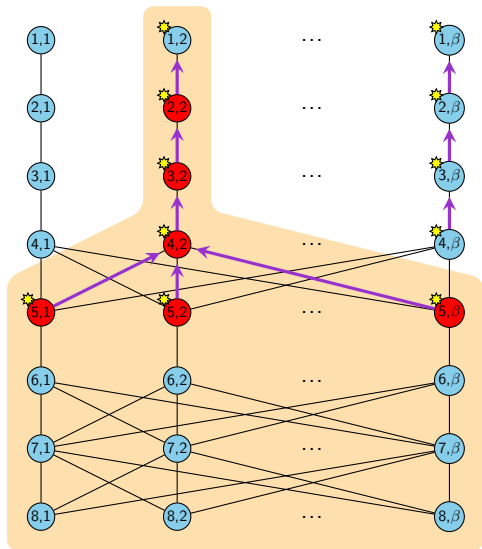
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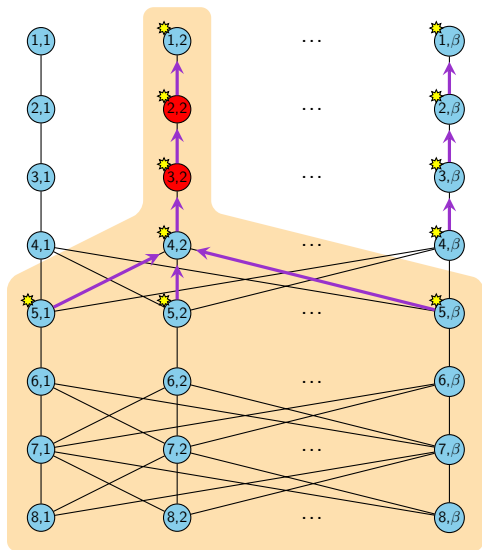
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$\beta^3$

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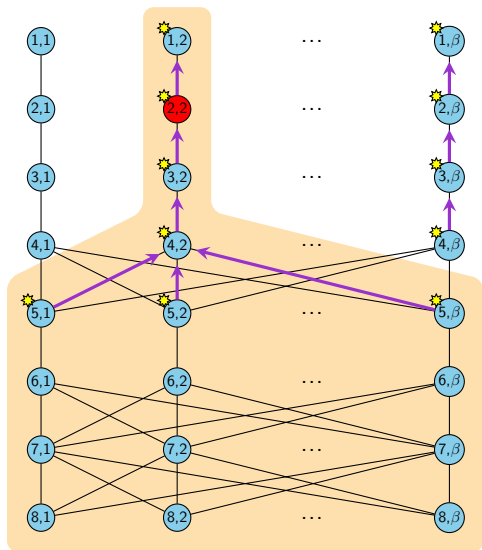
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Key:

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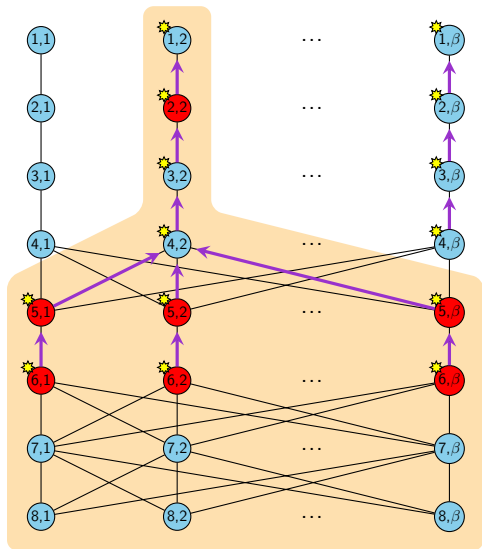
# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$\beta^3$

Key:

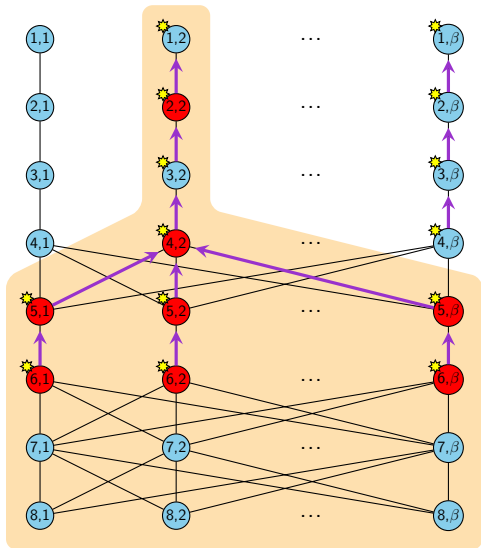
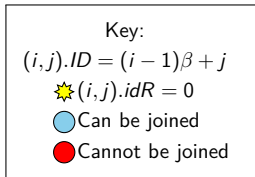
- $(i, j).ID = (i - 1)\beta + j$
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Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

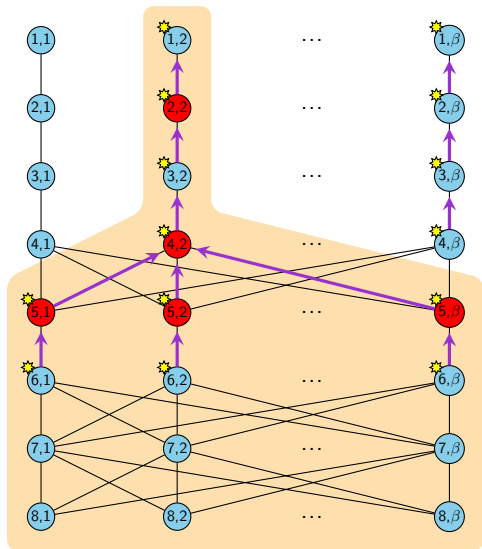
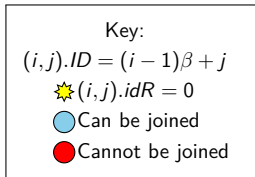
$\beta^3$



# Datta et al, 2011

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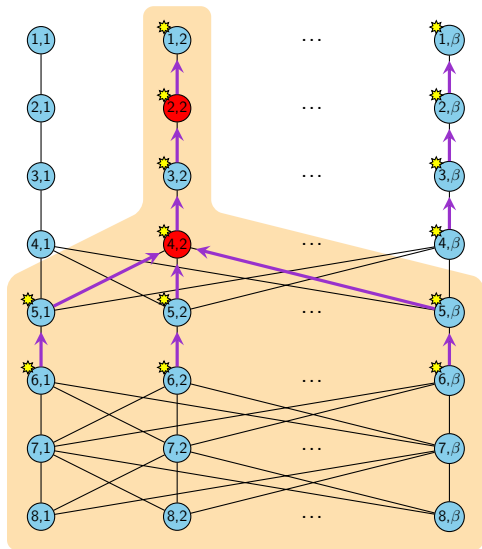
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$\beta^3$

Key:

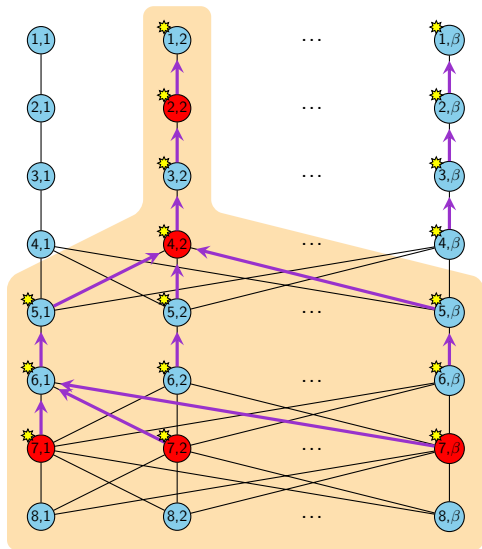
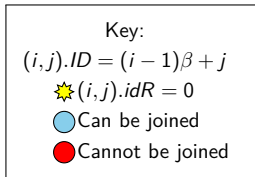
- $(i, j).ID = (i - 1)\beta + j$
- ★  $(i, j).idR = 0$
- Can be joined
- Cannot be joined



# Datta et al, 2011

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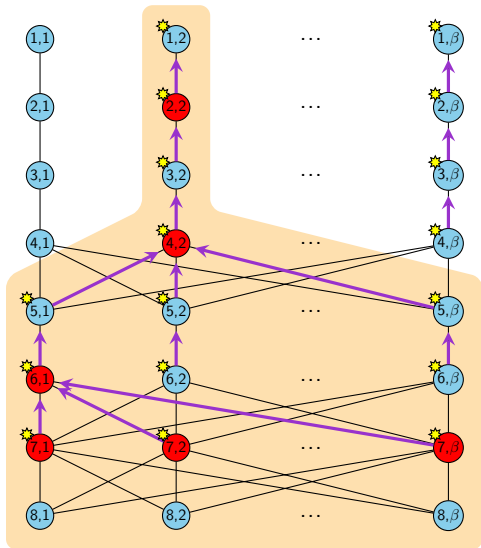
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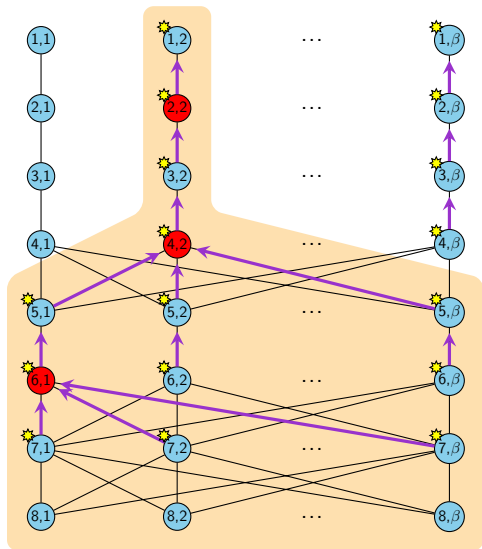
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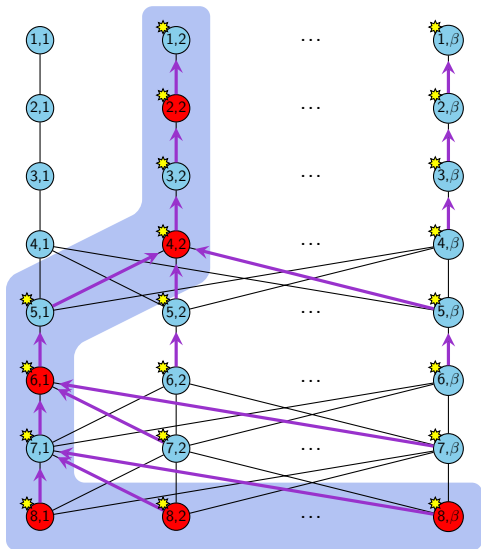
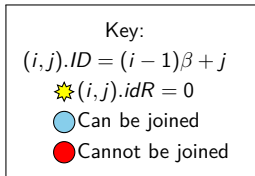




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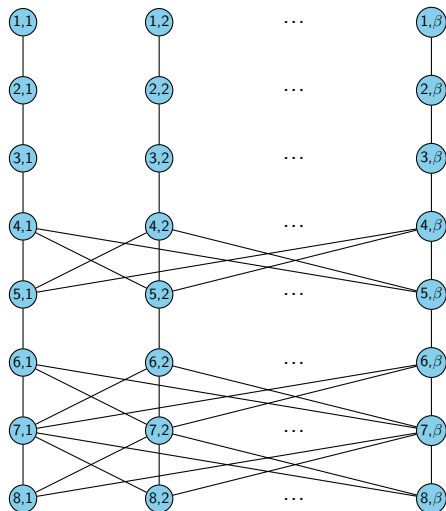
Key:

$(i, j).ID = (i - 1)\beta + j$

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● Cannot be joined



# Datta et al, 2011

Execution in  $\Omega(n^4)$  steps:  $\beta = \frac{n}{8}$

$$\beta = \Omega(n) \Rightarrow \beta^4 = \Omega(n^4)$$

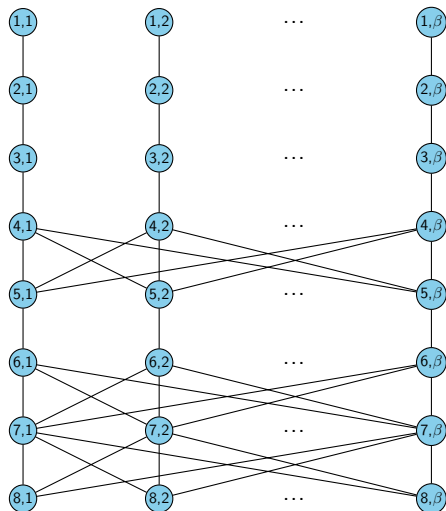
Key:

$$(i, j).ID = (i - 1)\beta + j$$

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● Can be joined

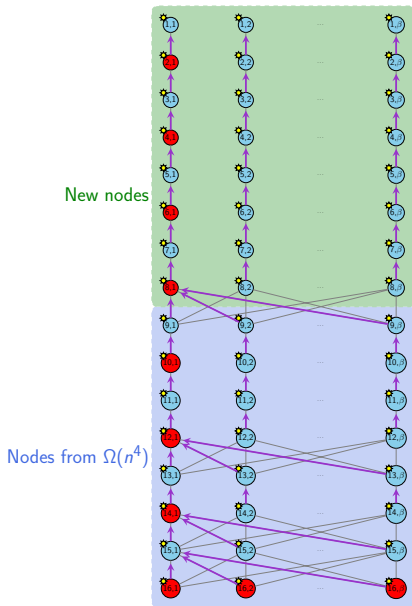
● Cannot be joined



# Datta et al, 2011

Network for  $\Omega(n^5)$  steps

$\forall \alpha \geq 3, \exists$  networks and executions in  $\Omega(n^{\alpha+1})$  steps.



# Perspectives

## Goal

Design a self-stabilizing leader election algorithm that stabilizes in  $O(\mathcal{D})$  rounds.

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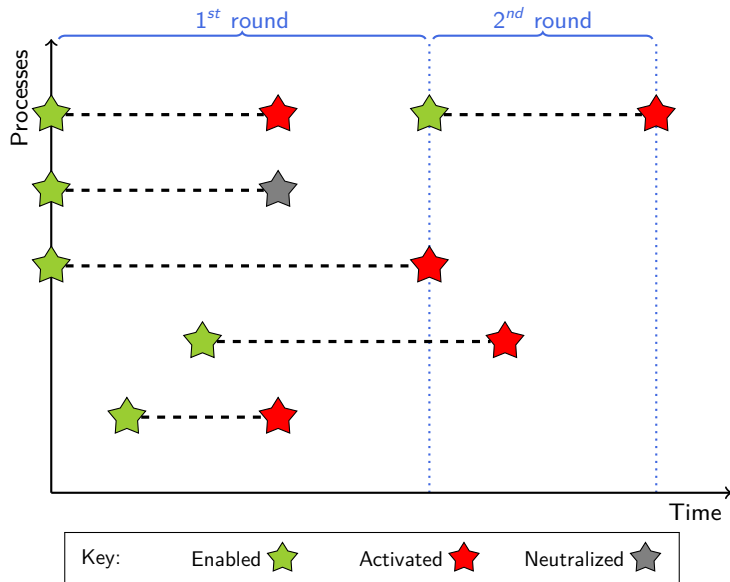
- Unfair daemon
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- Without any global knowledge : ??

Thank you for your attention.

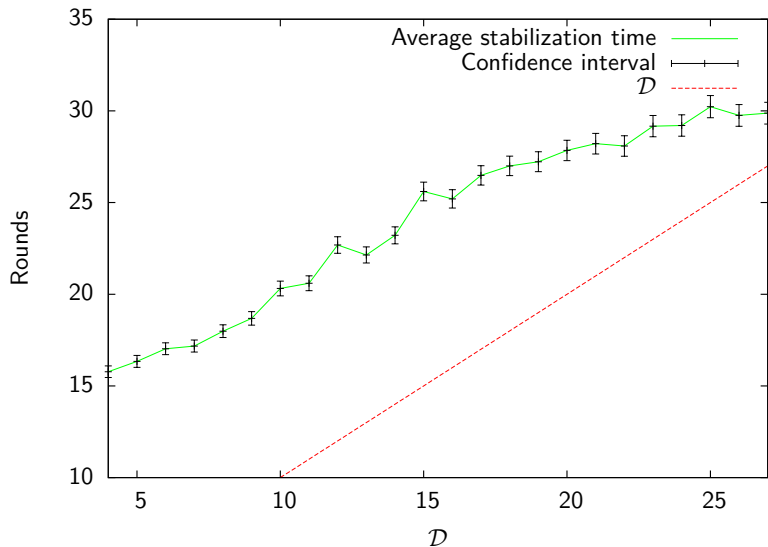
Do you have any questions ?



# Rounds

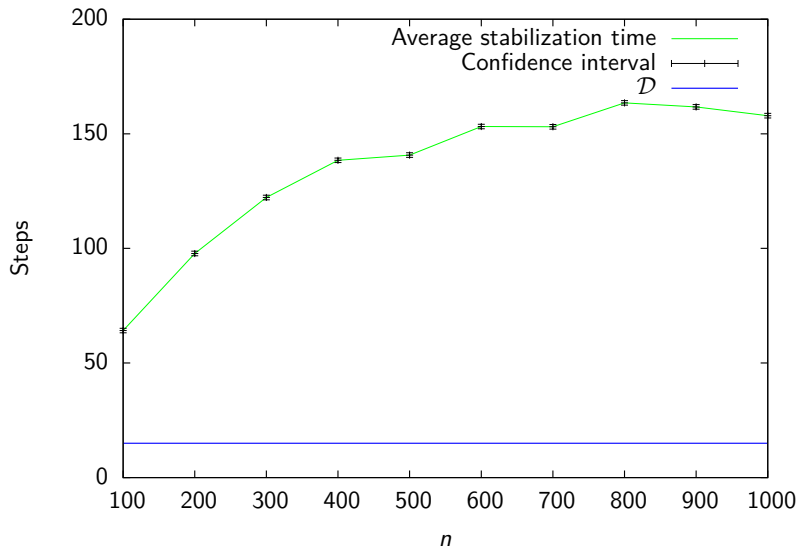


# Experimental Results



Average stabilization time in rounds in UDGs ( $n = 1000$ )

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Average stabilization time in steps in UDGs ( $\mathcal{D} = 15$ )